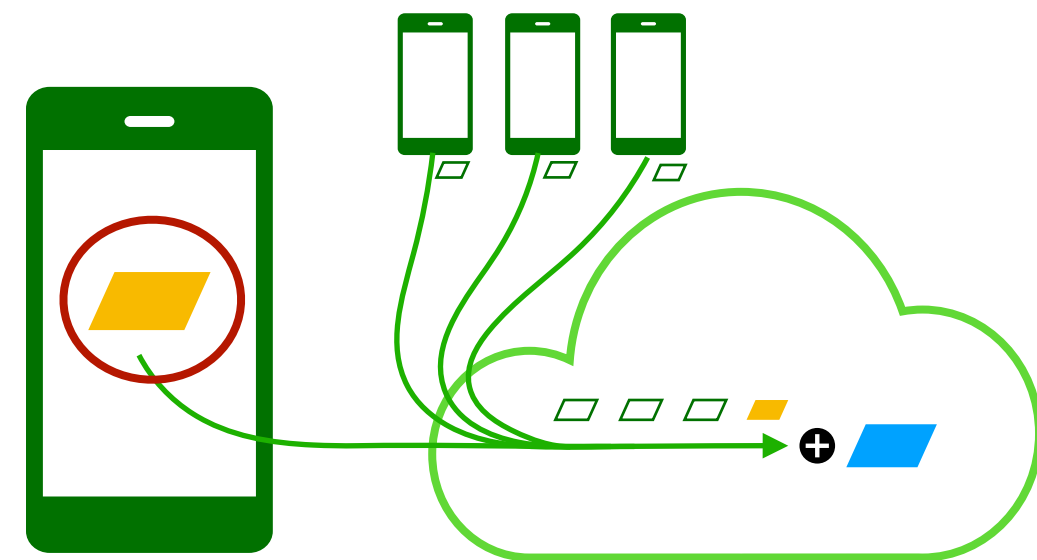


Lotto: Secure Participant Selection against Adversarial Servers in Federated Learning

Zhifeng Jiang, Peng Ye, Shiqi He, Wei Wang, Ruichuan Chen, Bo Li



Private learning on the edge

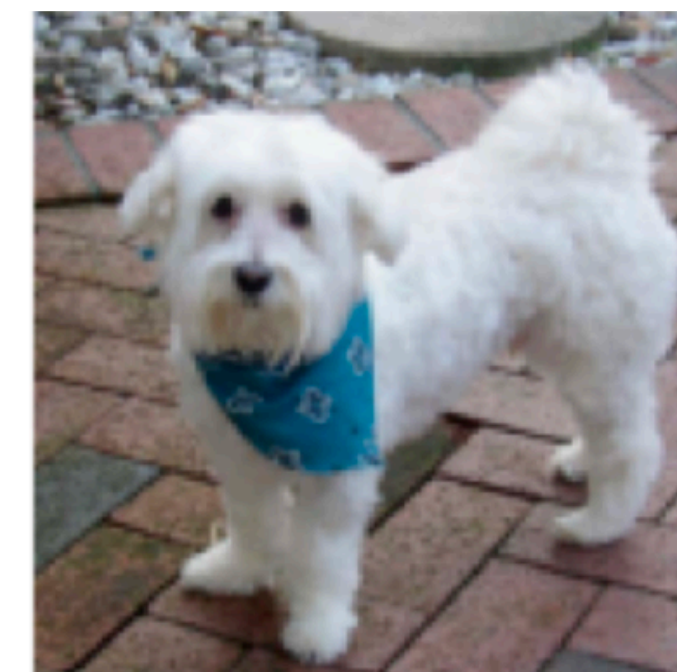


**Privacy-Enhancing
Technique**

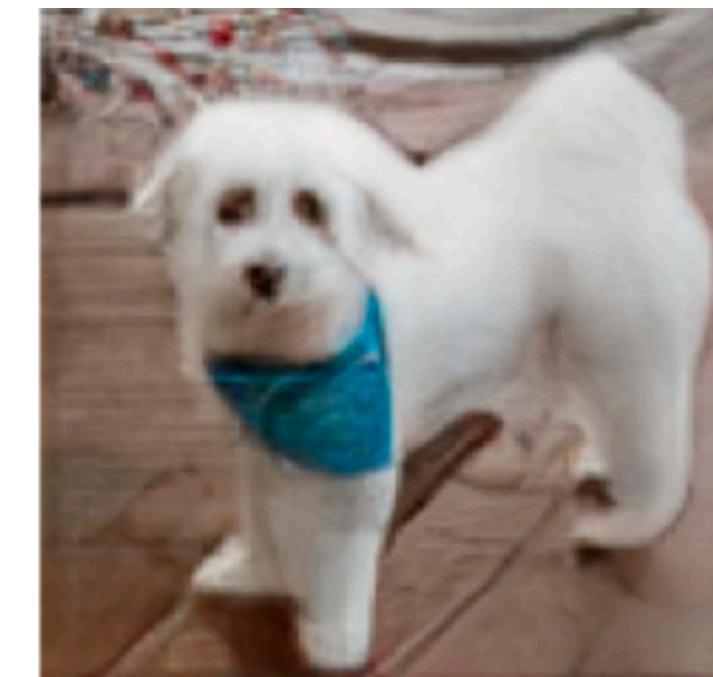
Federated Learning¹

Privacy Guarantee

Data kept on premises



Ground truth



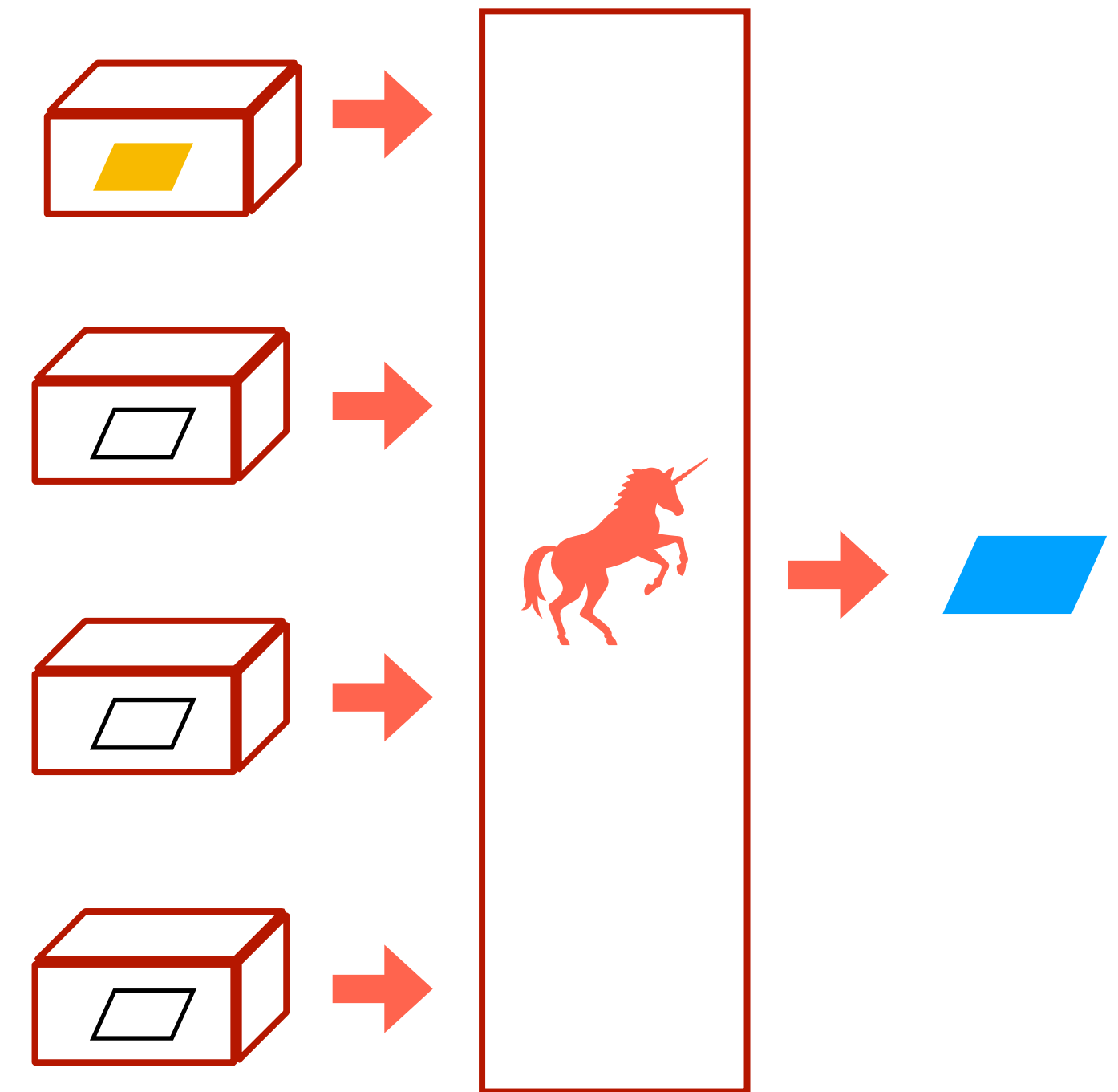
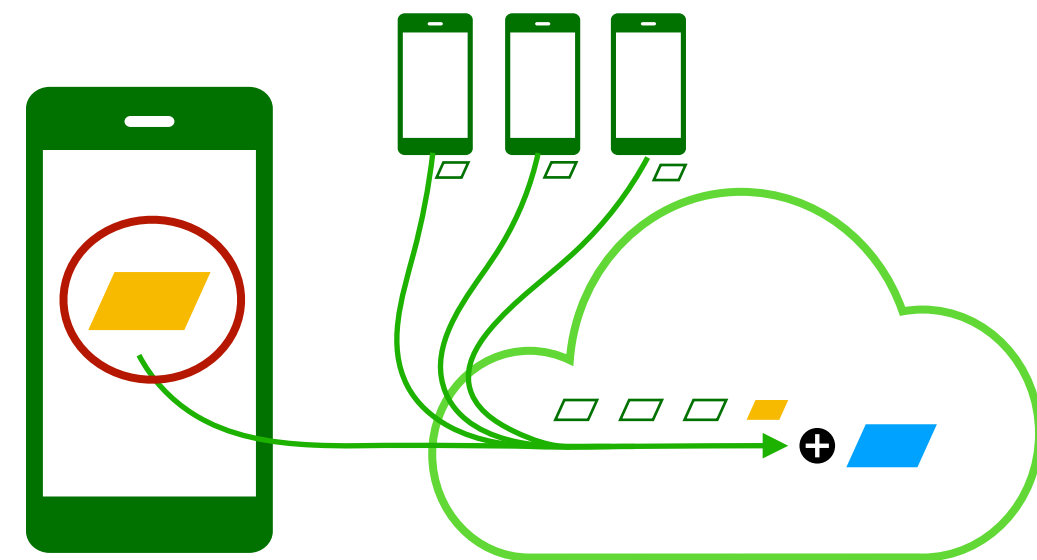
Reconstructed

Problem: Data can be reconstructed
from **local model updates**²

¹McMahan et al. "Communication-Efficient Learning of Deep Networks from Decentralized Data", In AISTATS '17

²Yue et al. "Gradient Obfuscation Gives a False Sense of Security in Federated Learning", In Security '23

Private learning on the edge



Privacy-Enhancing Technique	Federated Learning ¹	Secure Aggregation ^{3,4}
Privacy Guarantee	Data kept on premises	Local updates unseen

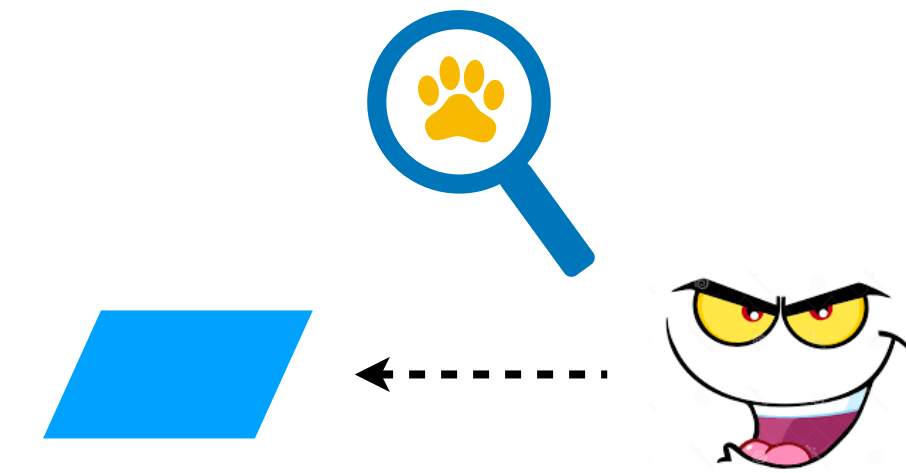
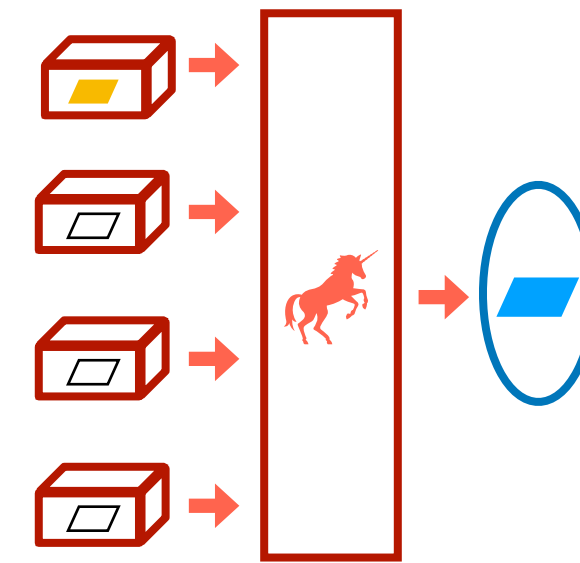
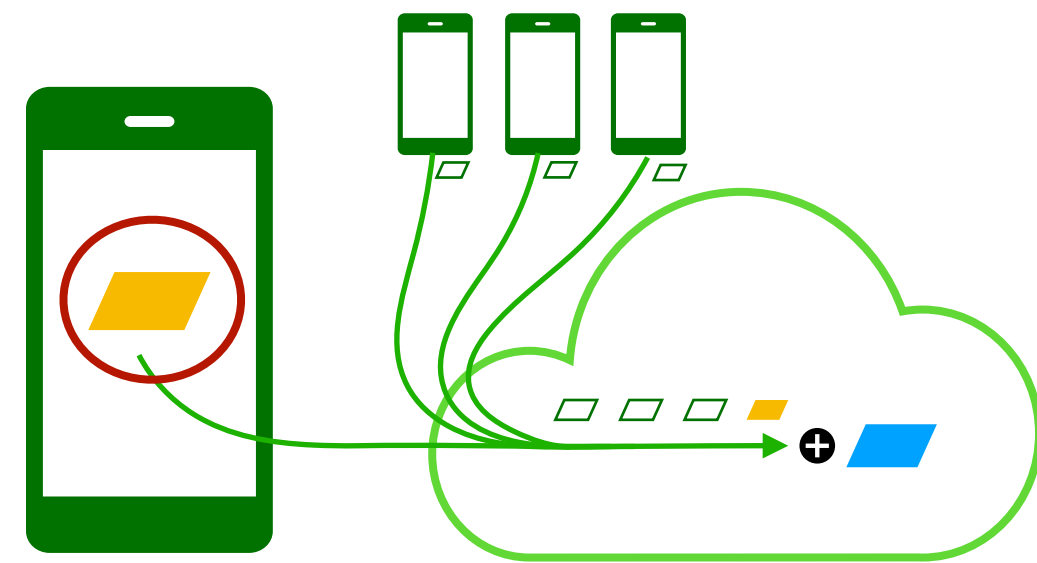
¹McMahan et al. "Communication-Efficient Learning of Deep Networks from Decentralized Data", In AISTATS '17

²Yue et al. "Gradient Obfuscation Gives a False Sense of Security in Federated Learning", In Security '23

³Bonawitz et al. "Practical Secure Aggregation for Privacy-Preserving Machine Learning", In CCS '17

⁴Bell et al. "Secure Single-Server Aggregation with (Poly) Logarithmic Overhead", In CCS '20

Private learning on the edge



Privacy-Enhancing Technique	Federated Learning ¹	Secure Aggregation ^{3,4}
Privacy Guarantee	Data kept on premises	Local updates unseen

Problem: Data still has footprints in **global model update**⁵

¹McMahan et al. "Communication-Efficient Learning of Deep Networks from Decentralized Data", In AISTATS '17

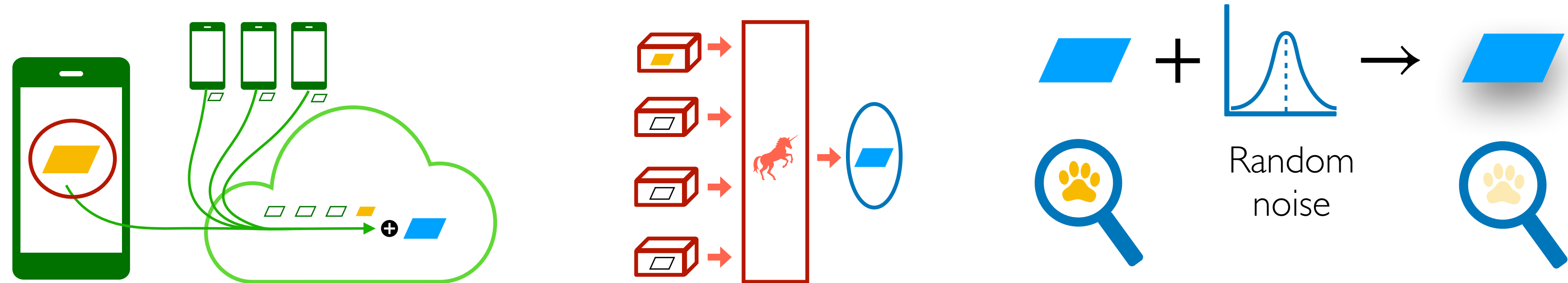
²Yue et al. "Gradient Obfuscation Gives a False Sense of Security in Federated Learning", In Security '23

³Bonawitz et al. "Practical Secure Aggregation for Privacy-Preserving Machine Learning", In CCS '17

⁴Bell et al. "Secure Single-Server Aggregation with (Poly) Logarithmic Overhead", In CCS '20

⁵Nasr et al. "Comprehensive Privacy Analysis of Deep Learning: Passive and Active White-box Inference Attacks against Centralized and Federated Learning", In S&P '19

Private learning on the edge



Privacy-Enhancing Technique	Federated Learning ¹	Secure Aggregation ^{3,4}	Differential Privacy ⁶
Privacy Guarantee	Data kept on premises	Local updates unseen	Global update leaks little about any client

¹McMahan et al. "Communication-Efficient Learning of Deep Networks from Decentralized Data", In AISTATS '17

²Yue et al. "Gradient Obfuscation Gives a False Sense of Security in Federated Learning", In Security '23

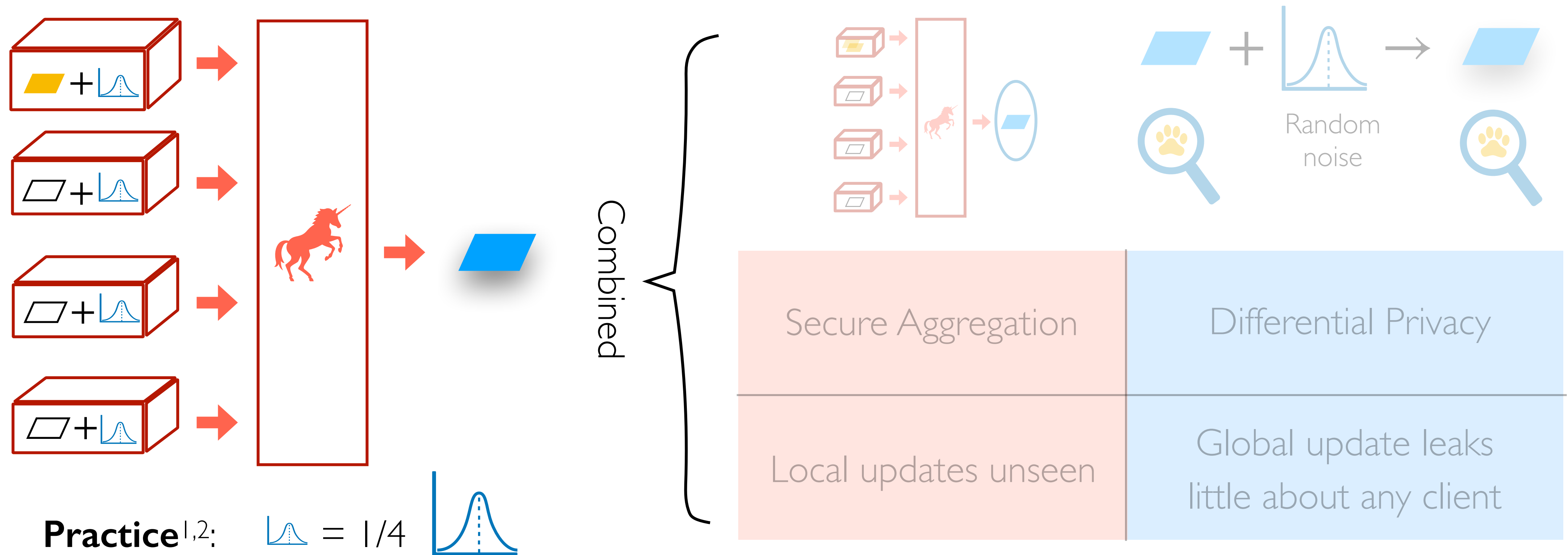
³Bonawitz et al. "Practical Secure Aggregation for Privacy-Preserving Machine Learning", In CCS '17

⁴Bell et al. "Secure Single-Server Aggregation with (Poly) Logarithmic Overhead", In CCS '20

⁵Nasr et al. "Comprehensive Privacy Analysis of Deep Learning: Passive and Active White-box Inference Attacks against Centralized and Federated Learning", In S&P '19

⁶Cynthia. "Differential Privacy", 06.

Private learning on the edge



Each client adds an **even share** of the target noise to its local model update

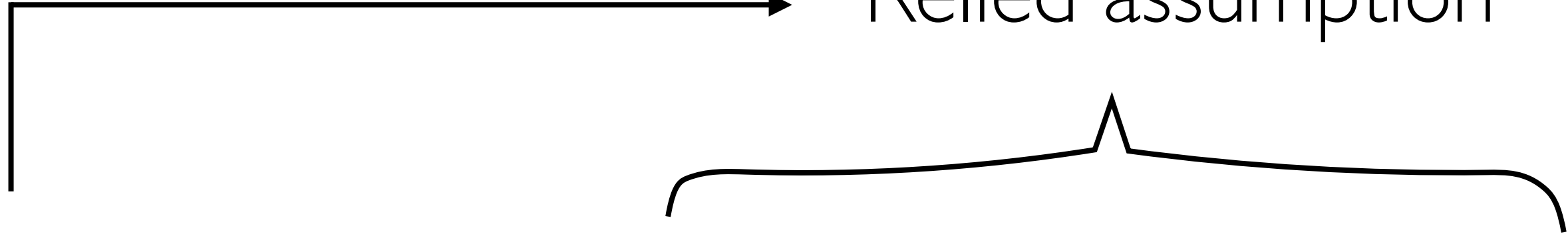
¹Kairouz et al. "The Distributed Discrete Gaussian Mechanism for Federated Learning with Secure Aggregation", In ICML '21

²Agarwal. "The Skellam Mechanism for Differentially Private Federated Learning", In NeurIPS '21

Private learning on the edge

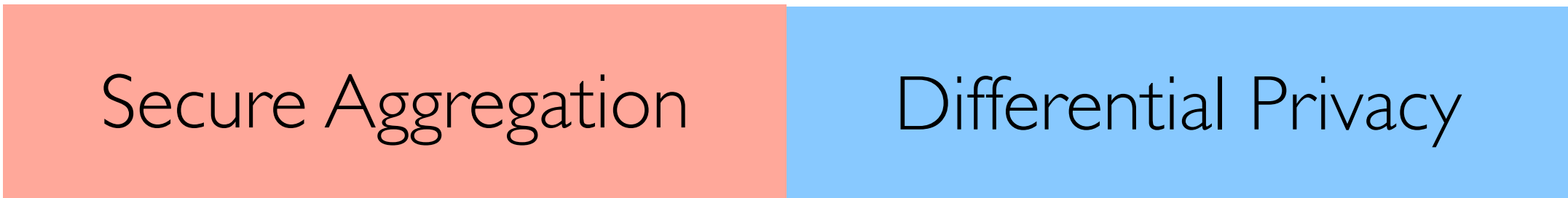
Privacy-Enhancing Technique	Federated Learning ¹	Secure Aggregation	Differential Privacy
Privacy Guarantee	Data kept on premises	Local updates unseen	Global update leaks little about any client

Need for Lotto

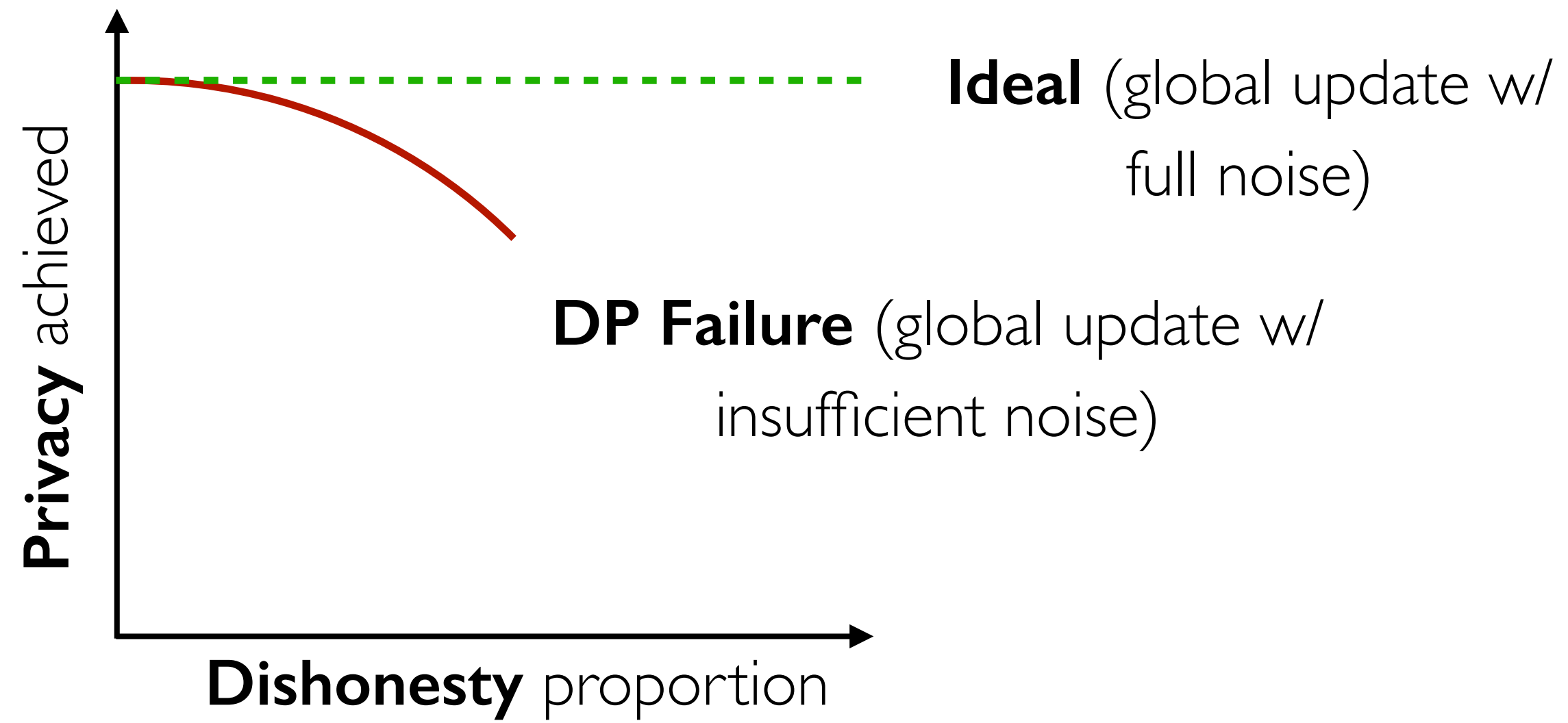
May **not** hold  Relied assumption

Privacy-Enhancing Technique	Federated Learning	Secure Aggregation	Differential Privacy
Privacy Guarantee	Data kept on premises	Local updates unseen	Global update leaks little about any client

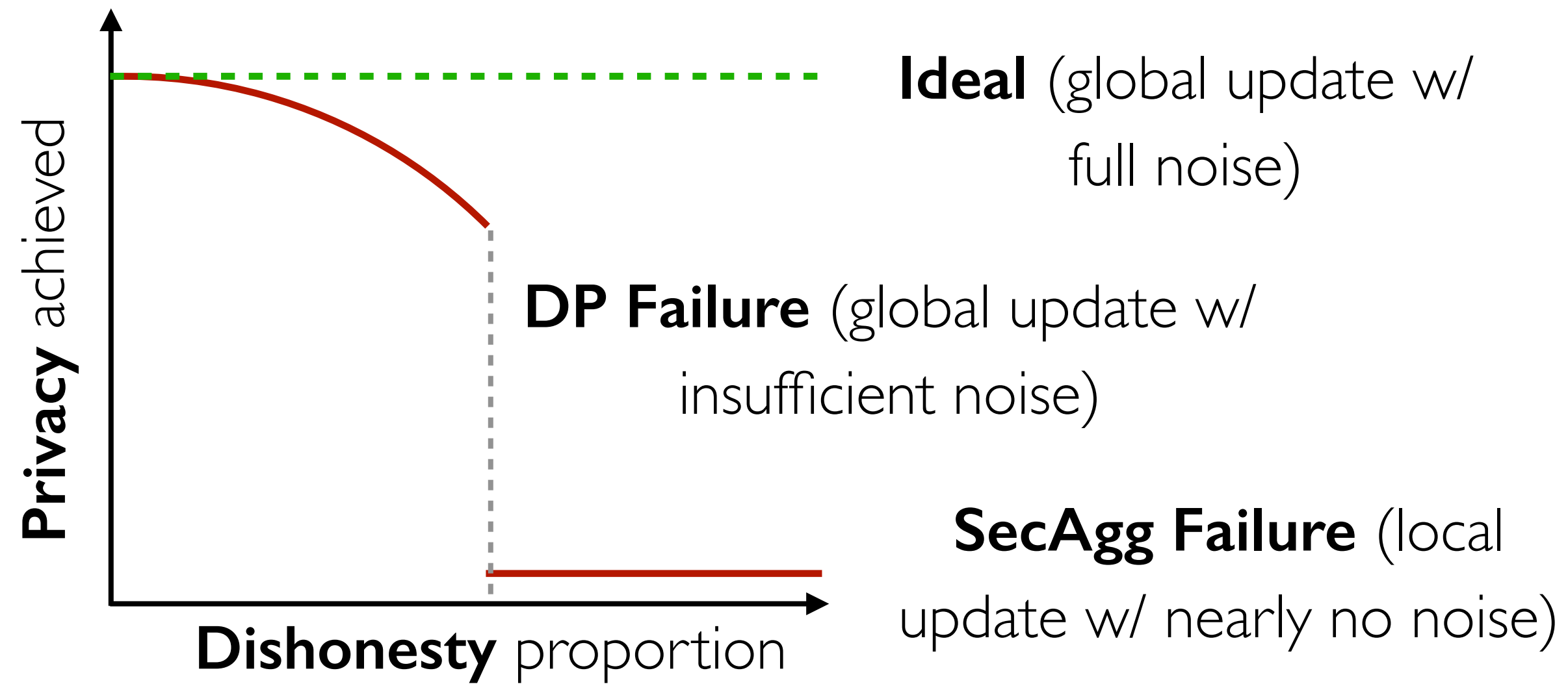
Need for Lotto



Need for Lotto



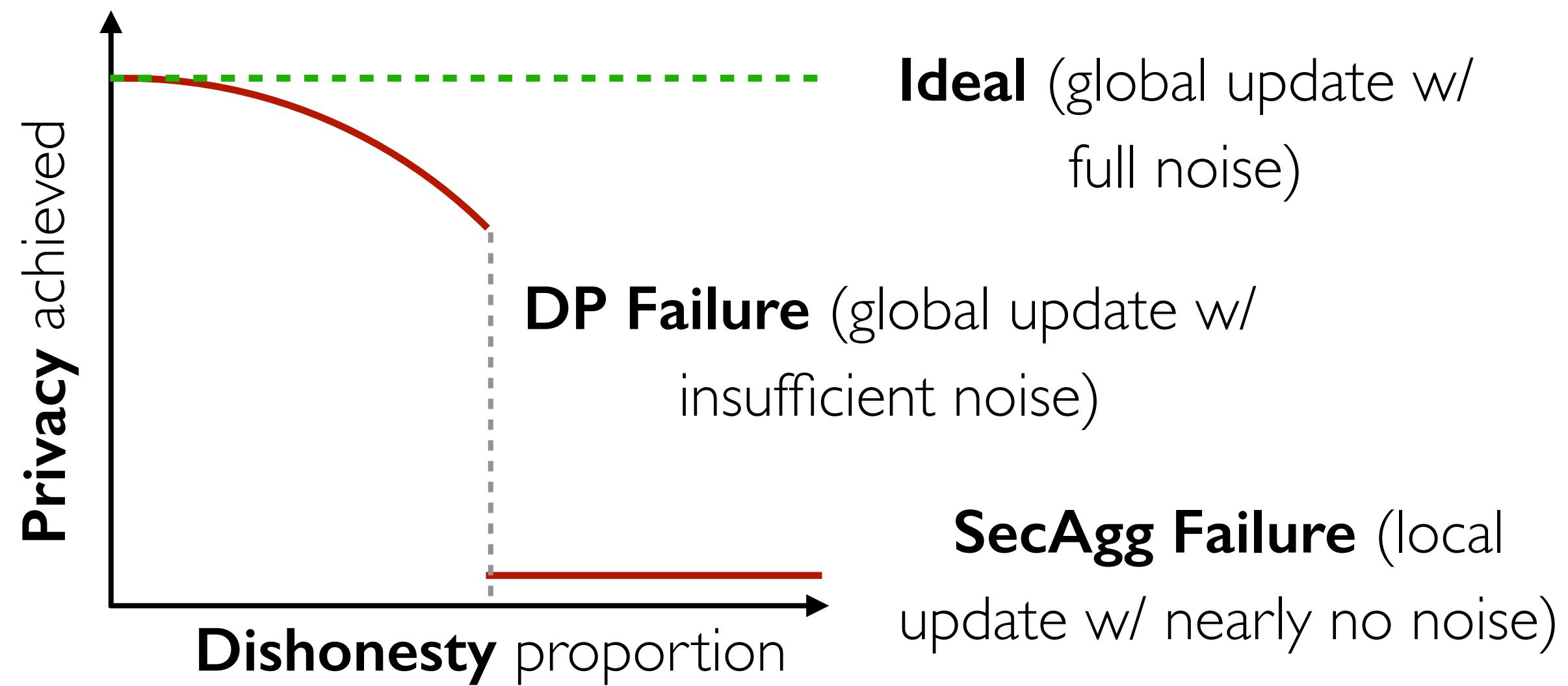
Need for Lotto



Secure Aggregation

Differential Privacy

Need for Lotto

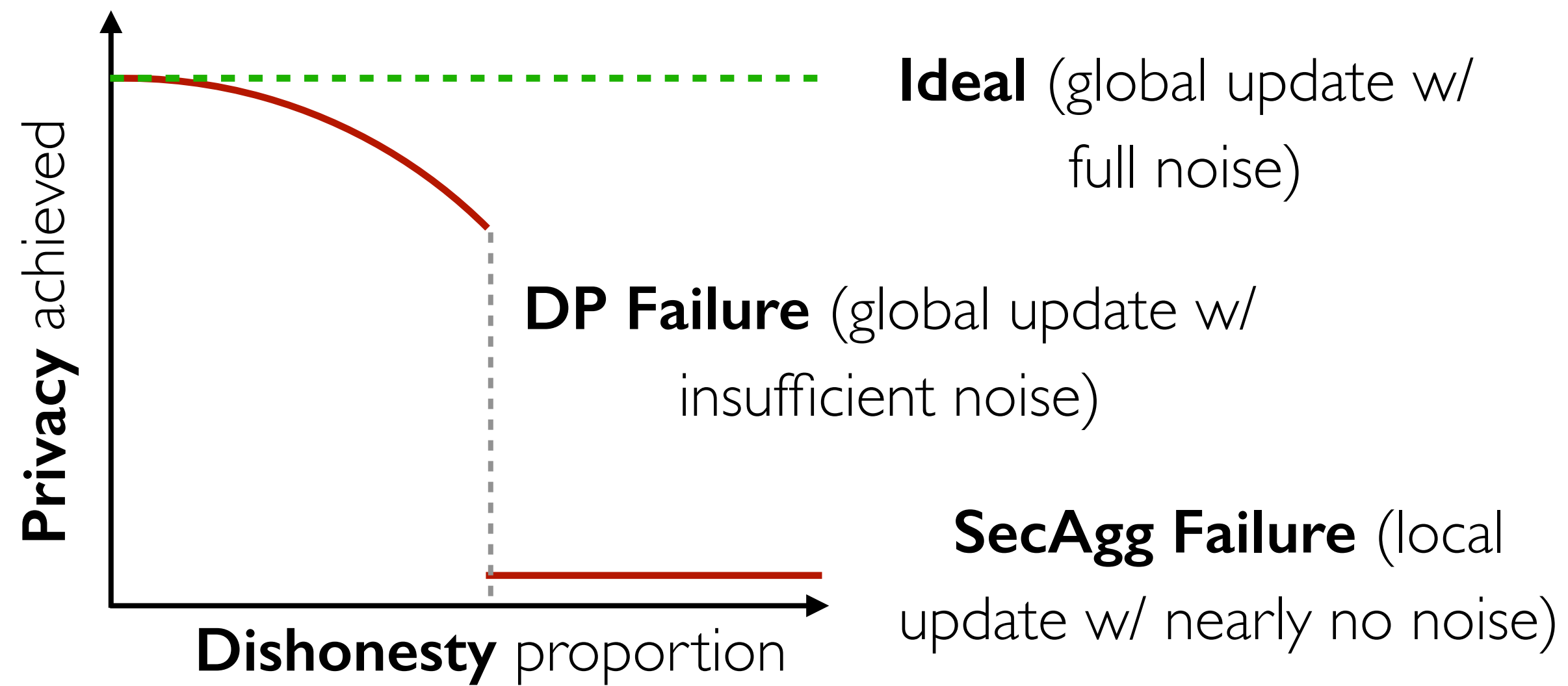


Assumption: honest participants

Secure Aggregation

Differential Privacy

Need for Lotto



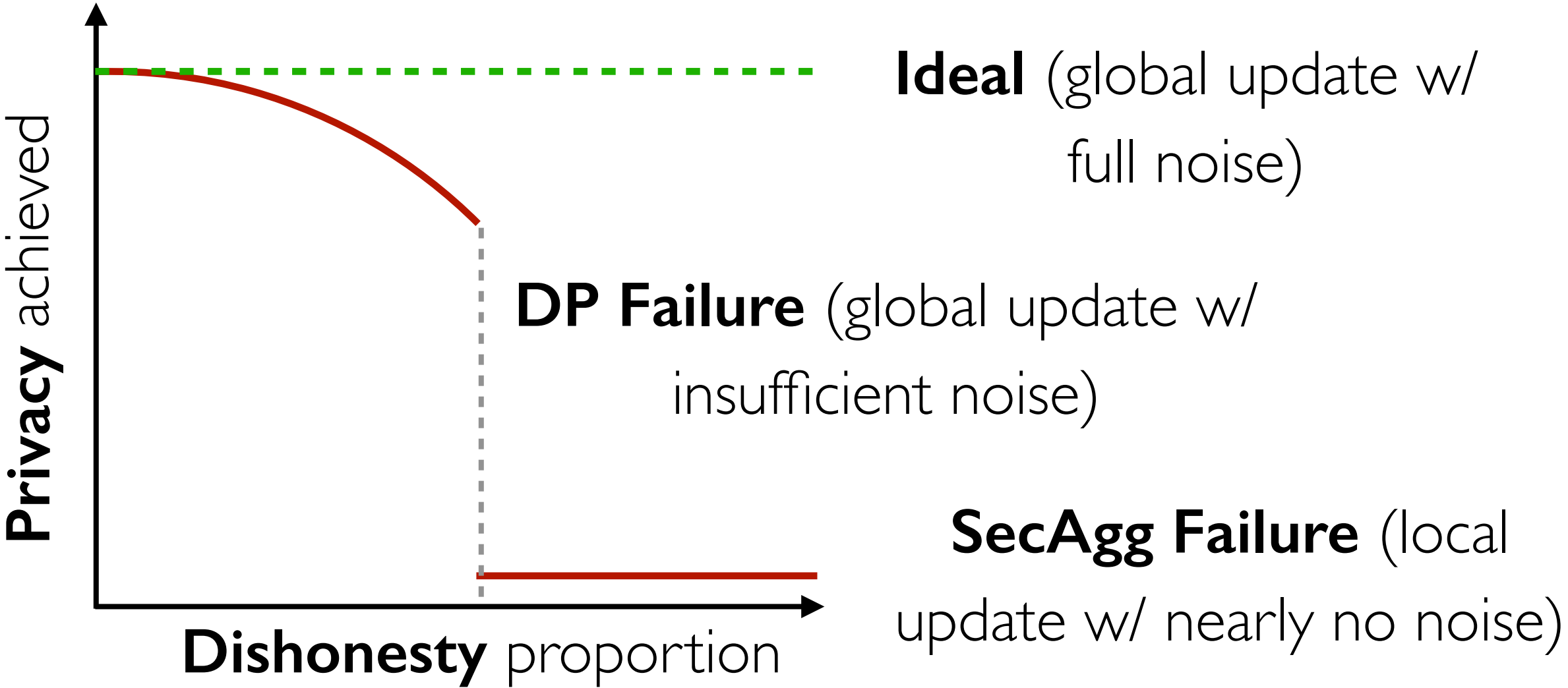
Assumption: honest participants

Secure Aggregation

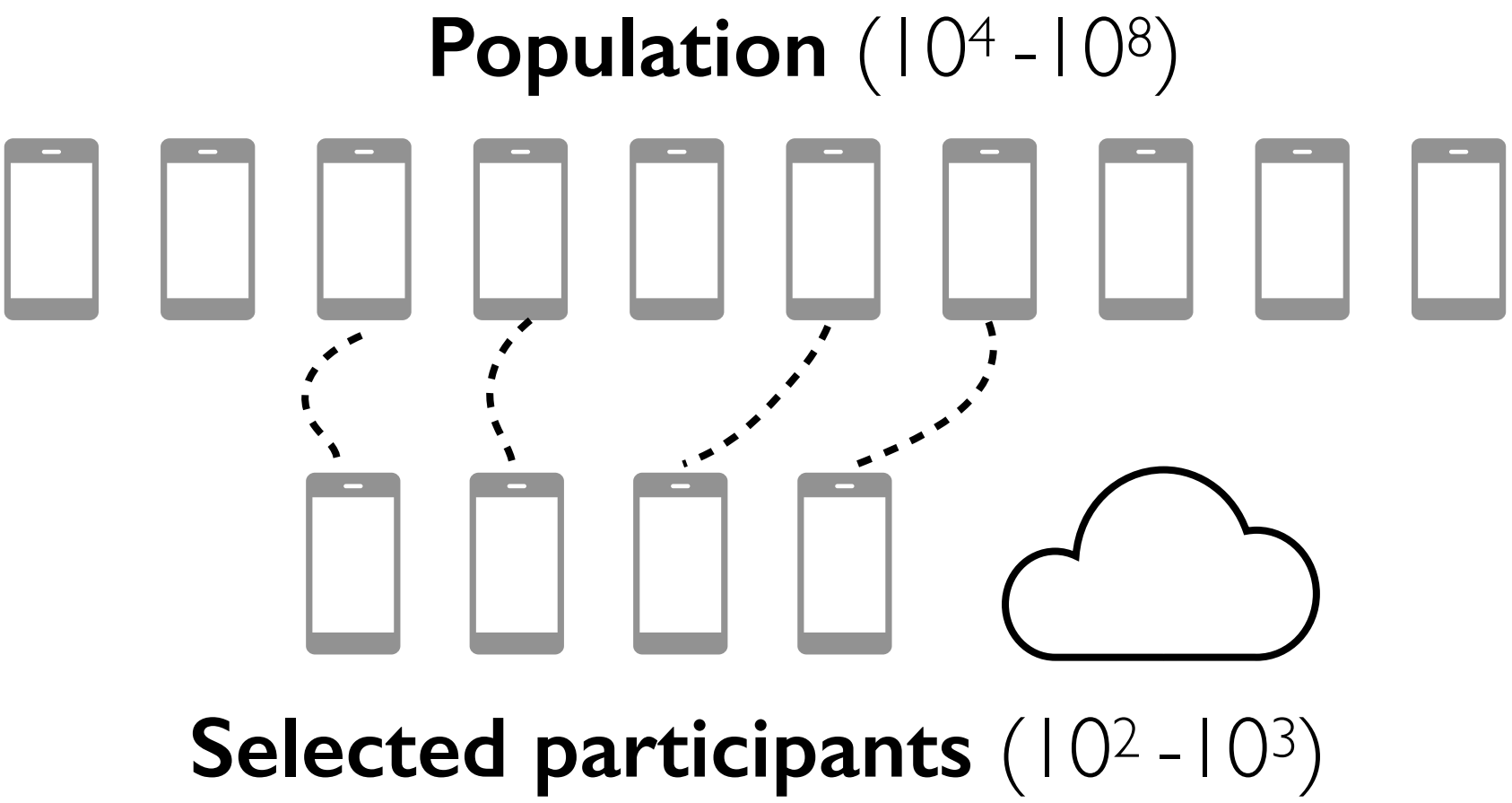
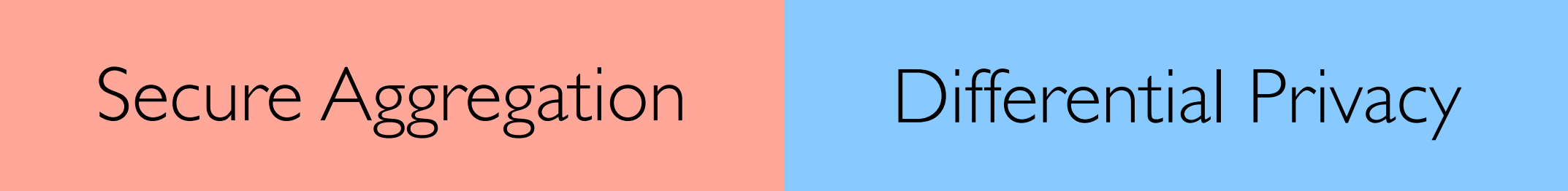
Differential Privacy

Federated Learning

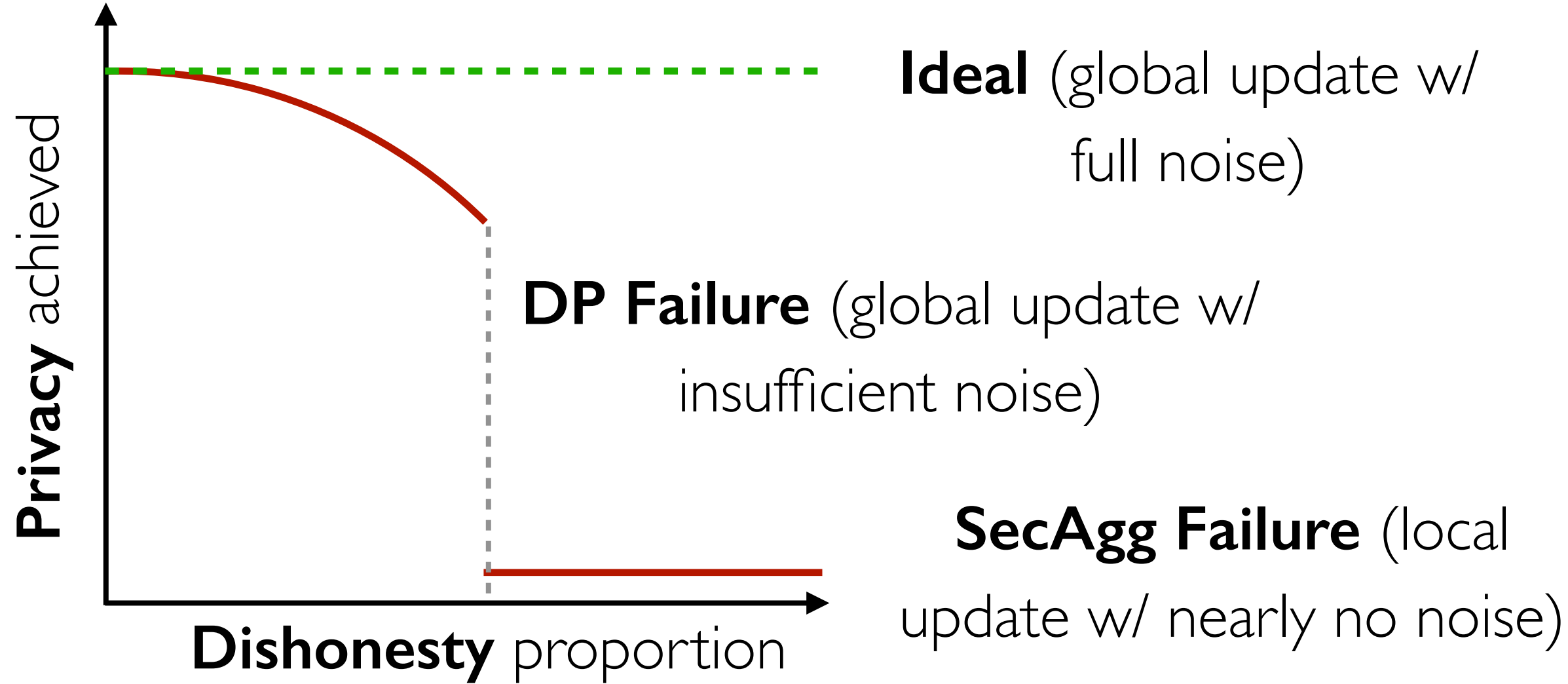
Need for Lotto



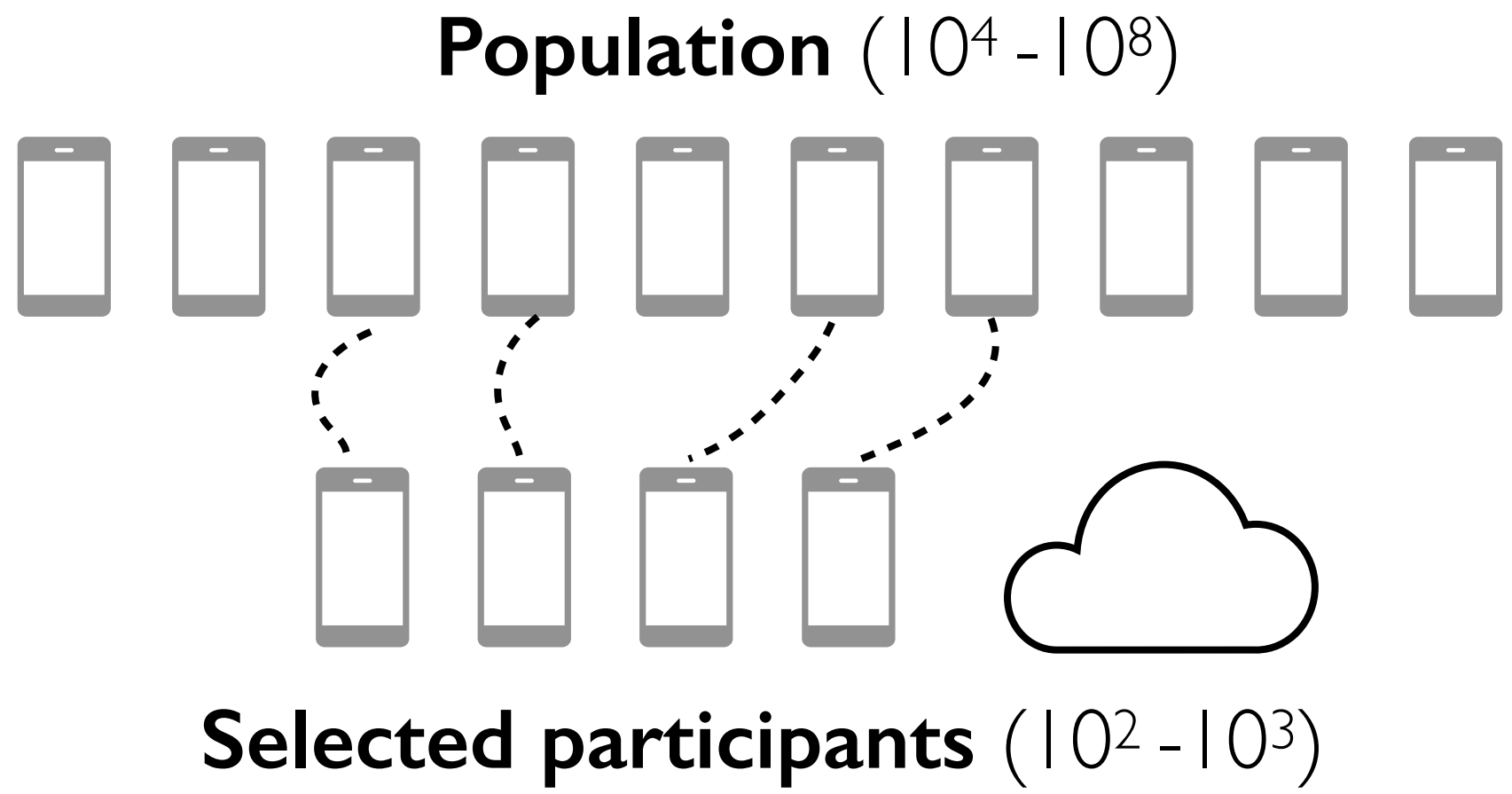
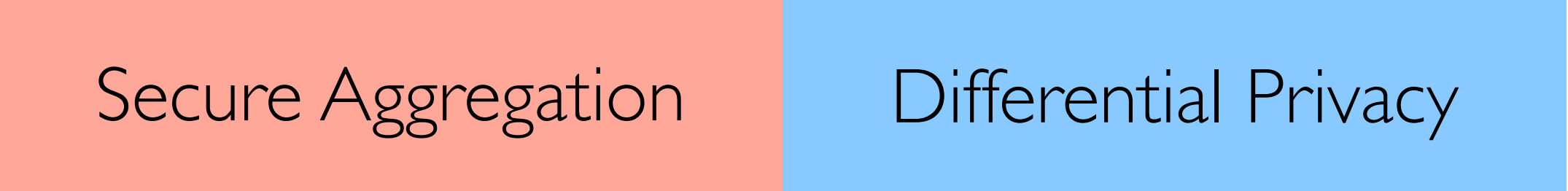
Assumption: honest participants



Need for Lotto



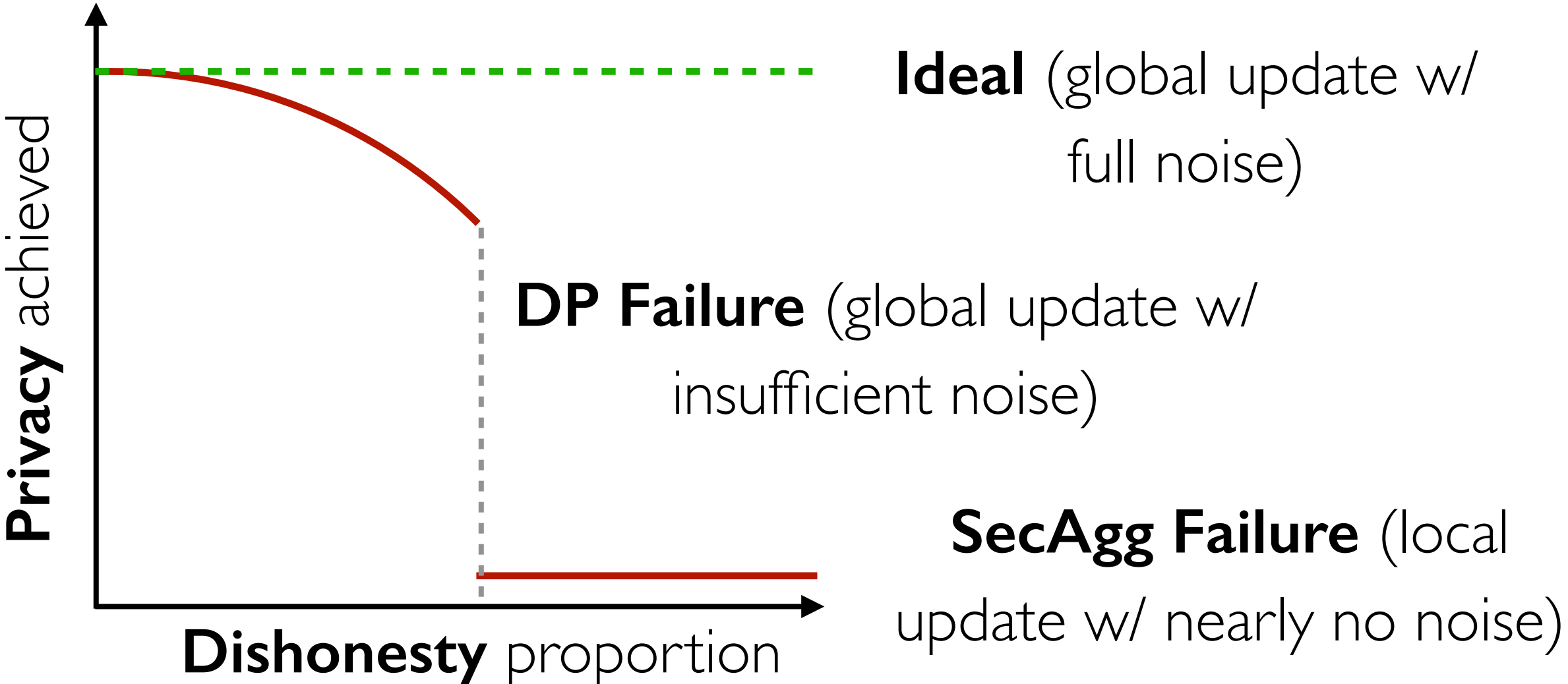
Assumption: honest participants



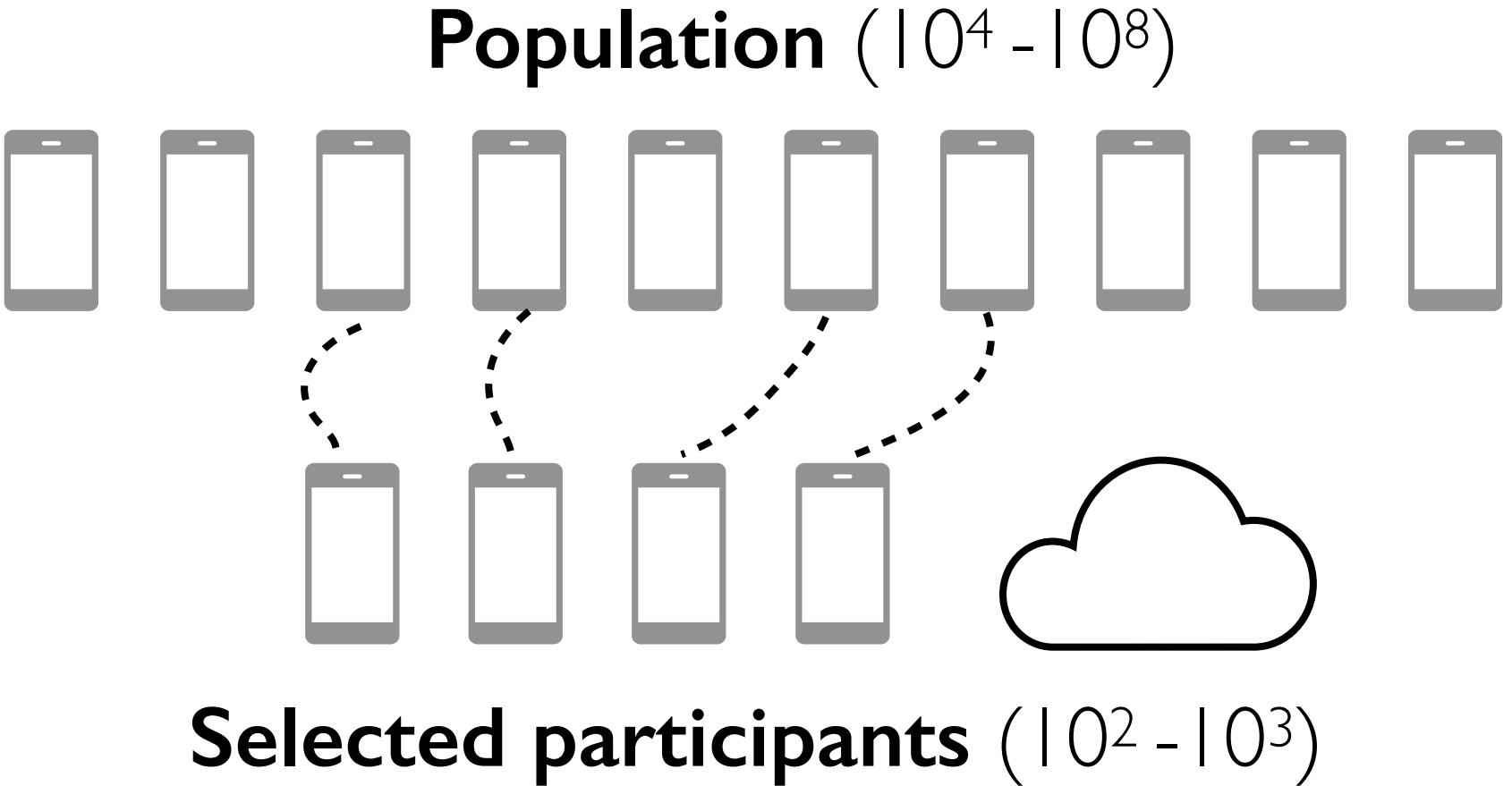
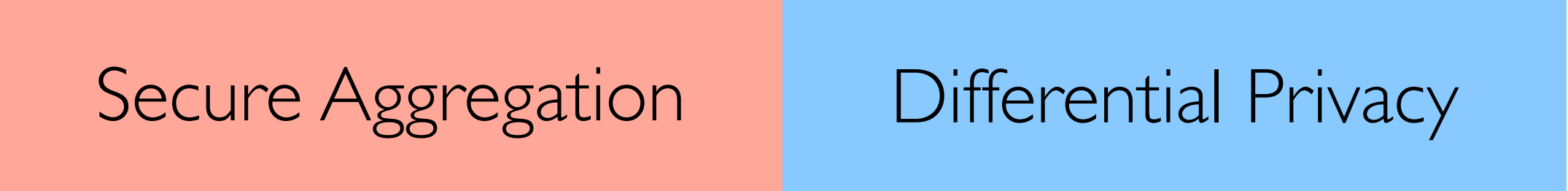
- **Random:** uniform chance



Need for Lotto



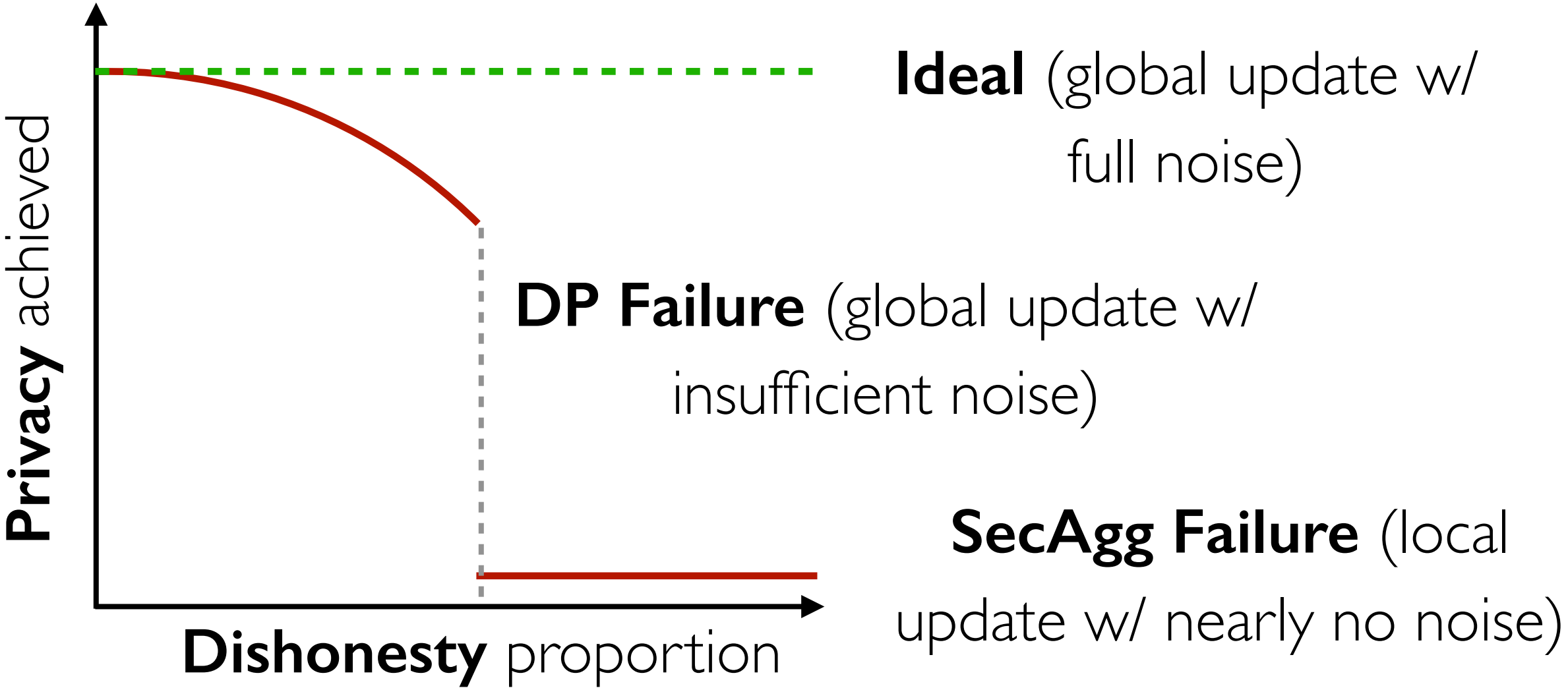
Assumption: honest participants



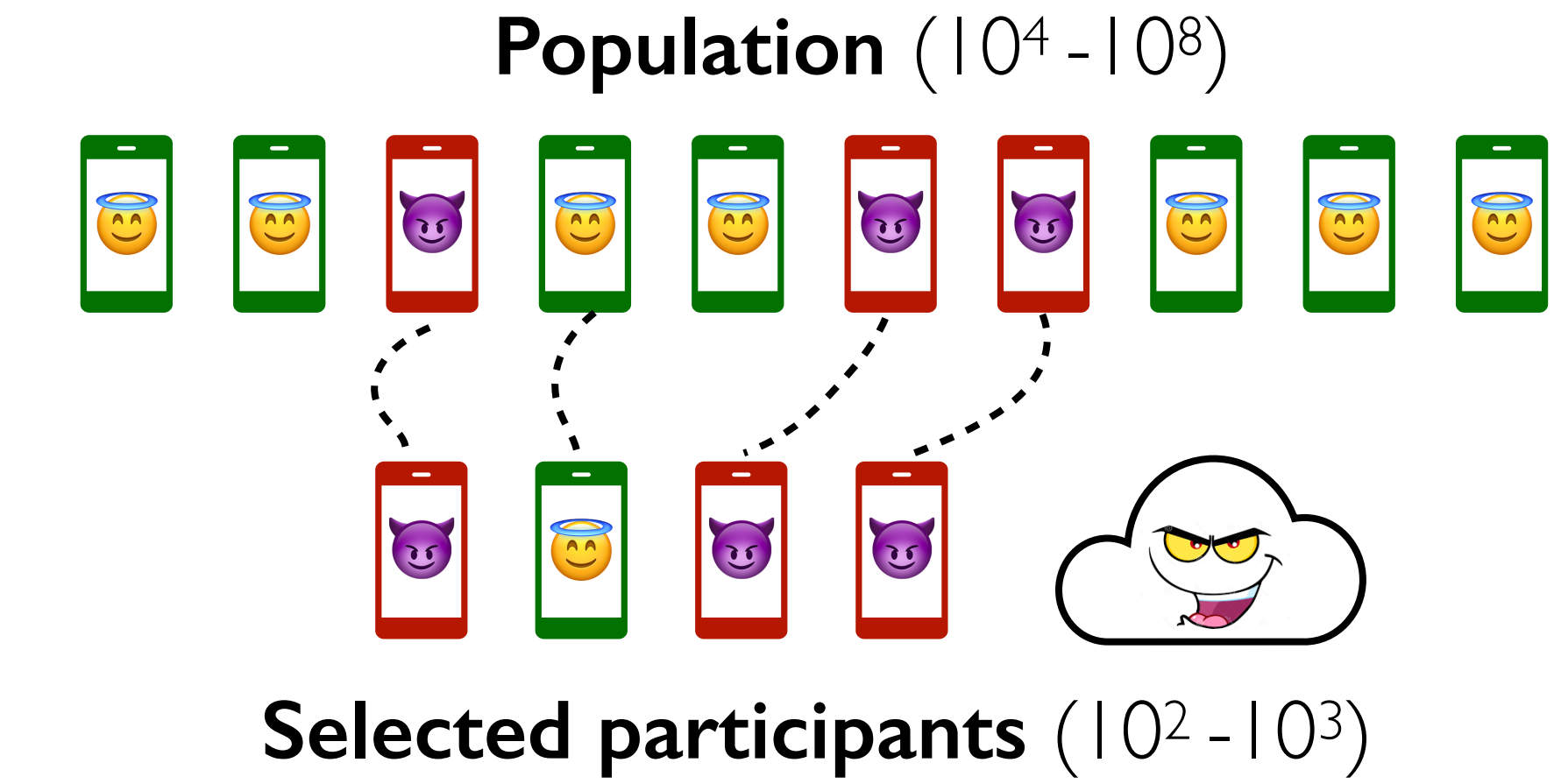
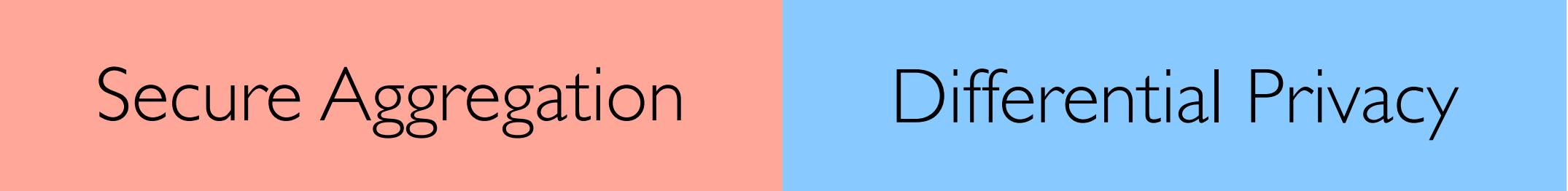
- **Random:** uniform chance
- **Informed:** “best-performing” clients are preferred (e.g., high speed and/or rich data)



Need for Lotto



Assumption: honest participants



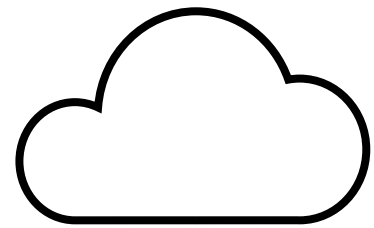
Problem: participant selection can be **manipulated** by the **malicious** server



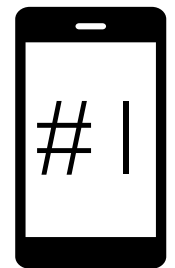
Lotto: Random selection

Lotto: Random selection

Current
round: 2



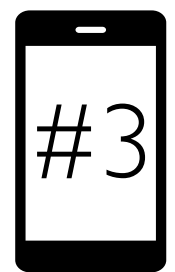
Randomness



$$\mathbf{RF}_{pk1}(2) = 9$$



$$\mathbf{RF}_{pk2}(2) = 1$$



$$\mathbf{RF}_{pk3}(2) = 7$$

...

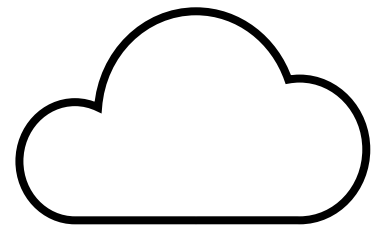
...

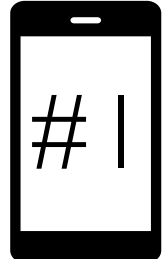

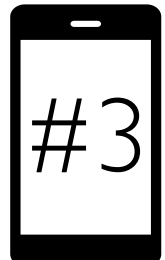
Public keys

Selection criteria: <3

Lotto: Random selection

Current
round: 2

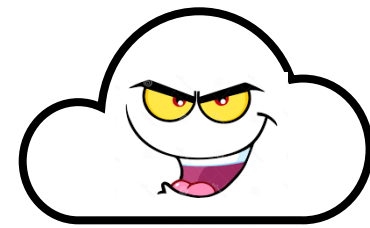
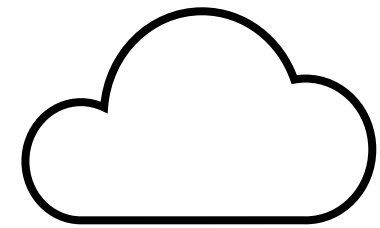


	Randomness	Select
 #1	$\mathbf{RF}_{pk1}(2) = 9$	No
 #2	$\mathbf{RF}_{pk2}(2) = 1$	Yes
 #3	$\mathbf{RF}_{pk3}(2) = 7$	No
...

Selection criteria: <3

Lotto: Random selection

Current round: 2



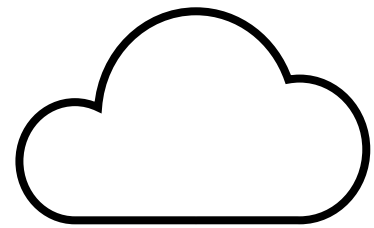
	Randomness	Select	Randomness	Select
#1	$\mathbf{RF}_{pk1}(2) = 9$	No		Yes
#2	$\mathbf{RF}_{pk2}(2) = 1$	Yes	Does NOT matter.	No
#3	$\mathbf{RF}_{pk3}(2) = 7$	No		No
...

Selection criteria: < 3

For dishonest majority

Lotto: Random selection

Current round: 2



	Randomness	Select	Randomness	Select
#1	$\mathbf{RF}_{pk1}(2) = 9$	No		Yes
#2	$\mathbf{RF}_{pk2}(2) = 1$	Yes	Does NOT matter.	No
#3	$\mathbf{RF}_{pk3}(2) = 7$	No		No
...

Selection criteria: < 3

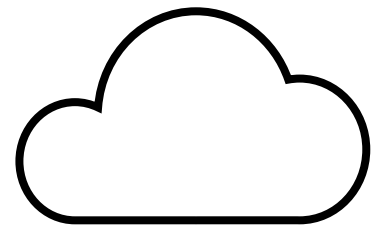
For dishonest majority

Potential approach:

- Mutual verification

Lotto: Random selection

Current round: 2



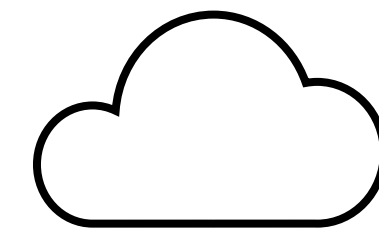
	Randomness	Select	Randomness	Select
#1	$RF_{pk1}(2) = 9$	No		Yes
#2	$RF_{pk2}(2) = 1$	Yes	Does NOT matter.	No
#3	$RF_{pk3}(2) = 7$	No		No
...

Selection criteria: < 3

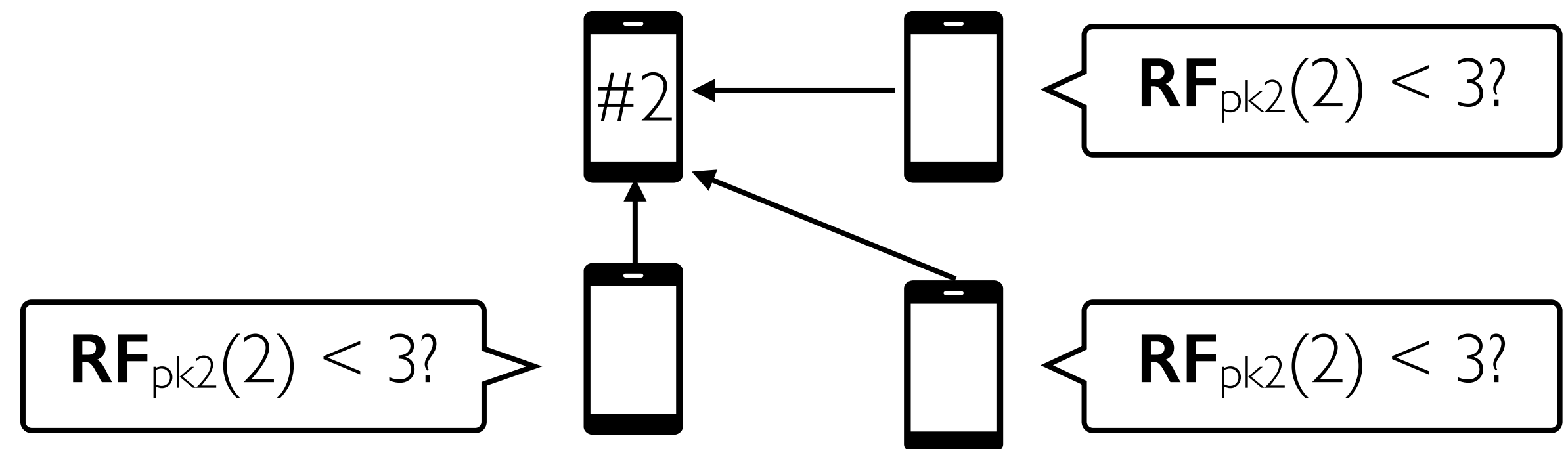
For dishonest majority

Potential approach:

- Mutual verification

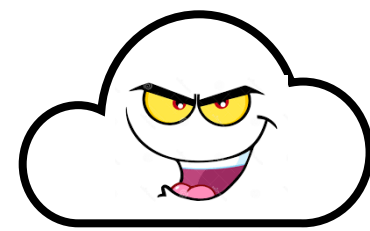
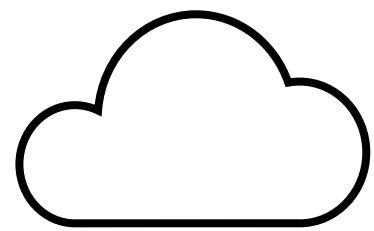


I select #2



Lotto: Random selection

Current round: 2



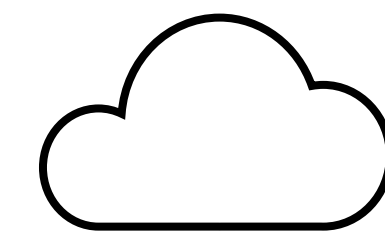
	Randomness	Select	Randomness	Select
#1	$RF_{pk1}(2) = 9$	No		Yes
#2	$RF_{pk2}(2) = 1$	Yes	Does NOT matter.	No
#3	$RF_{pk3}(2) = 7$	No		No
...

Selection criteria: <3

For dishonest majority

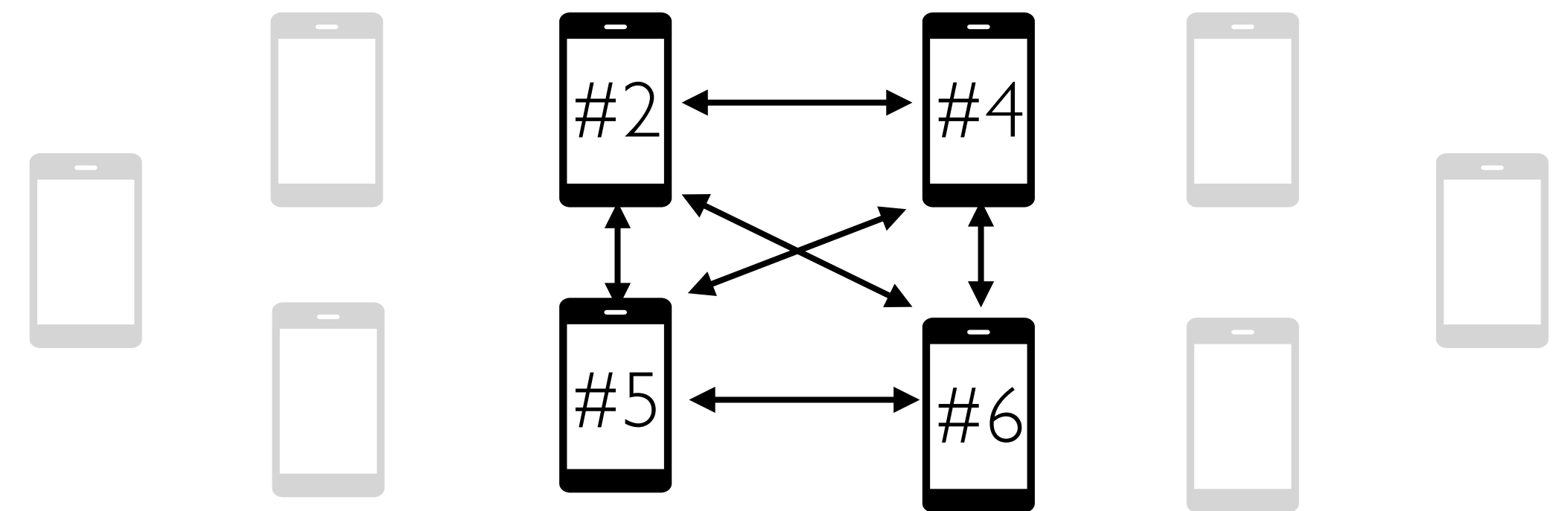
Potential approach:

- Mutual verification
- Only within participants ($10^2 - 10^3$)



I select #2, #4, #5, #6

Necessary →



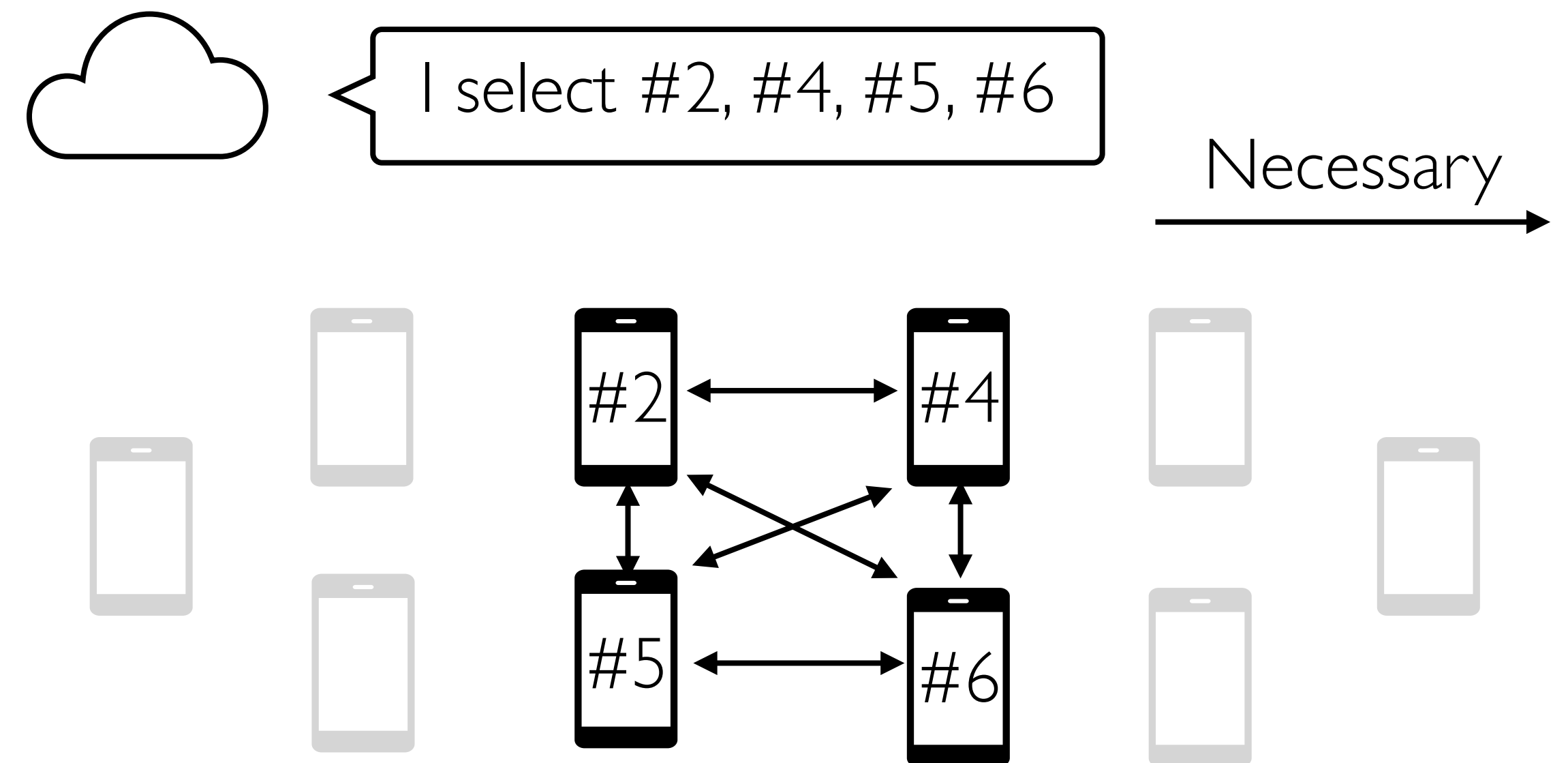
Lotto: Random selection

What is achieved:

Each participant
sees a list of peers

Potential approach:

- Mutual verification
- Only within participants ($10^2 - 10^3$)



Lotto: Random selection

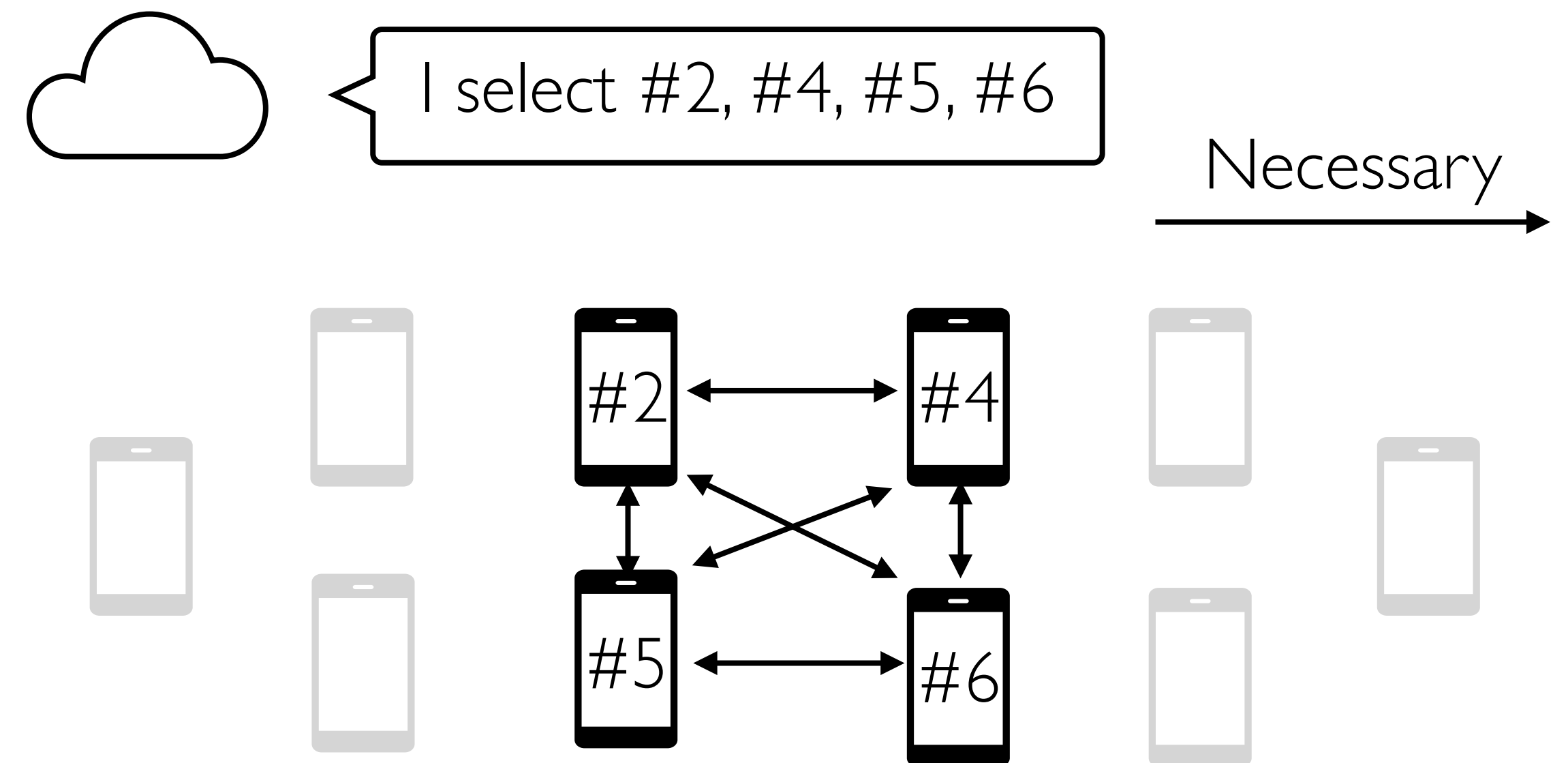
What is achieved:

Each participant
sees a list of peers who
presents only **by chance**.

E.g.,
$$\frac{\text{Selection criteria: } <3}{\text{Output range: } [0, 10)} = 3/10$$

Potential approach:

- Mutual verification
- Only within participants ($10^2 - 10^3$)



Lotto: Random selection

What is achieved:

Each participant

sees a list of peers who

presents only by chance.



What happens to the absent?

Lotto: Random selection

What is achieved:

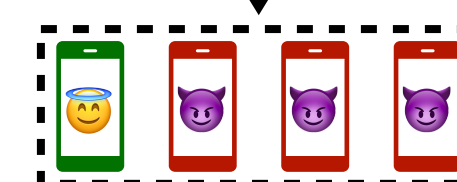
Each participant
sees a list of peers who
presents only by chance.



What happens to the absent?

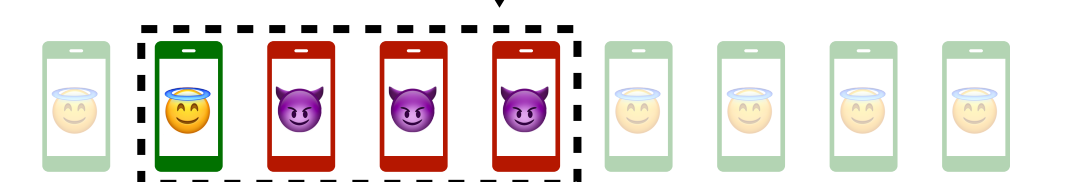
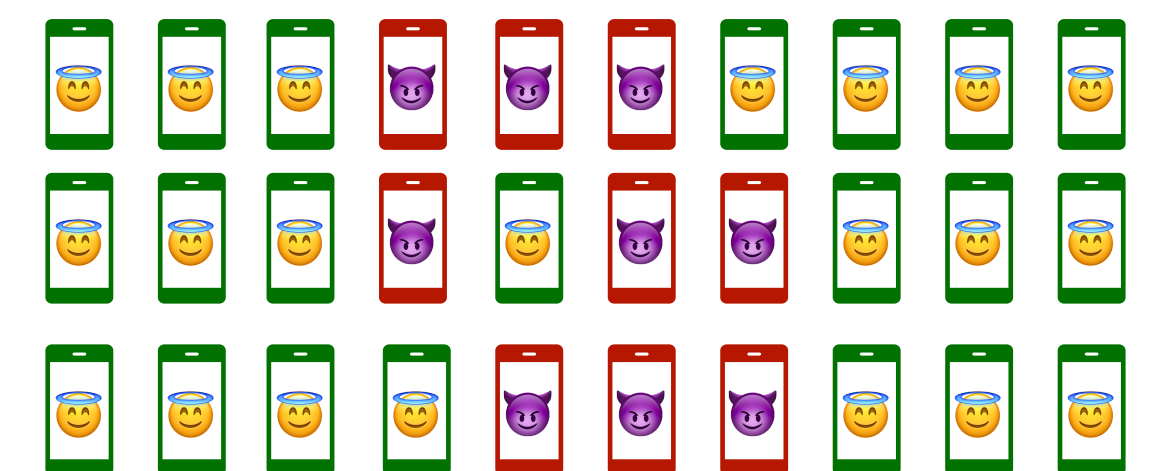
Problem: The server may arbitrarily
ignore honest clients

Ignore **before** selection



Selected

Ignore **after** selection



Lotto: Random selection

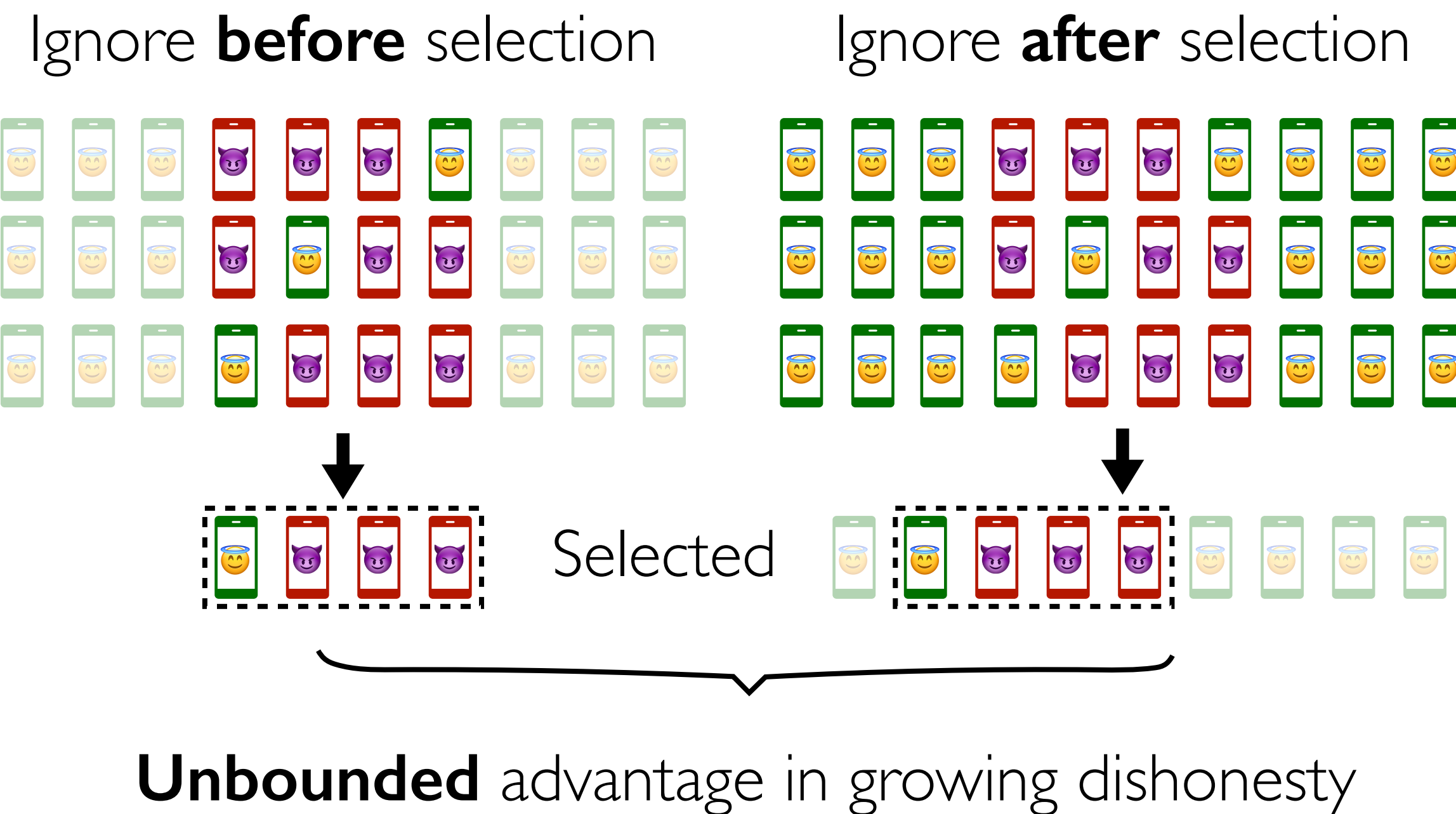
What is achieved:

Each participant
sees a list of peers who
presents only by chance.



What happens to the absent?

Problem: The server may arbitrarily
ignore honest clients



Lotto: Random selection

What is achieved:

Each participant
sees a list of peers who
presents only by chance.



What happens to the absent?

Solution: Enforce a **large enough list**
and a **small enough chance.**

Lotto: Random selection

What is achieved:

Each participant
sees a list of peers who
presents only by chance.



What happens to the absent?

Solution: Enforce a **large enough list**
and a **small enough chance.**

Example

- **len(list):** ≥ 200
- **Chance:** $\leq 0.1\%$

Lotto: Random selection

What is achieved:

Each participant
sees a list of peers who
presents only by chance.

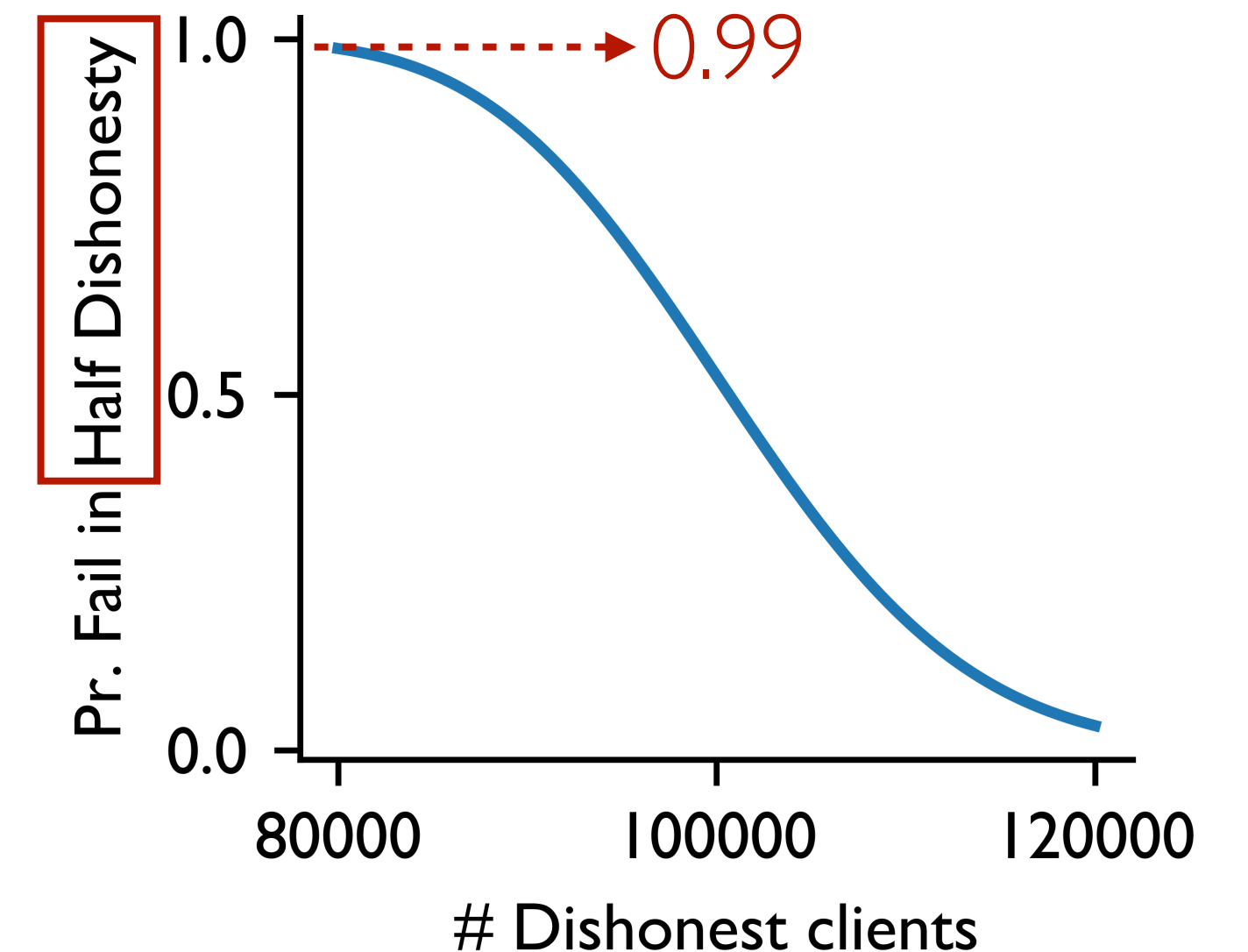


What happens to the absent?

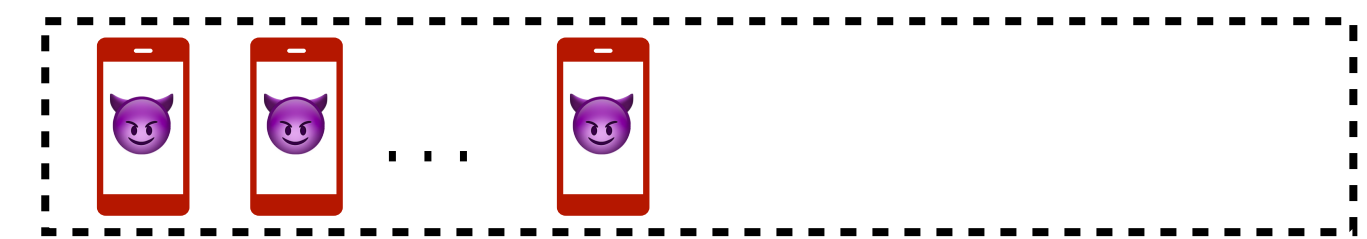
Solution: Enforce a **large enough list**
and a **small enough chance.**

Example

- **len(list):** ≥ 200
- **Chance:** $\leq 0.1\%$



Selected



Lotto: Random selection

What is achieved:

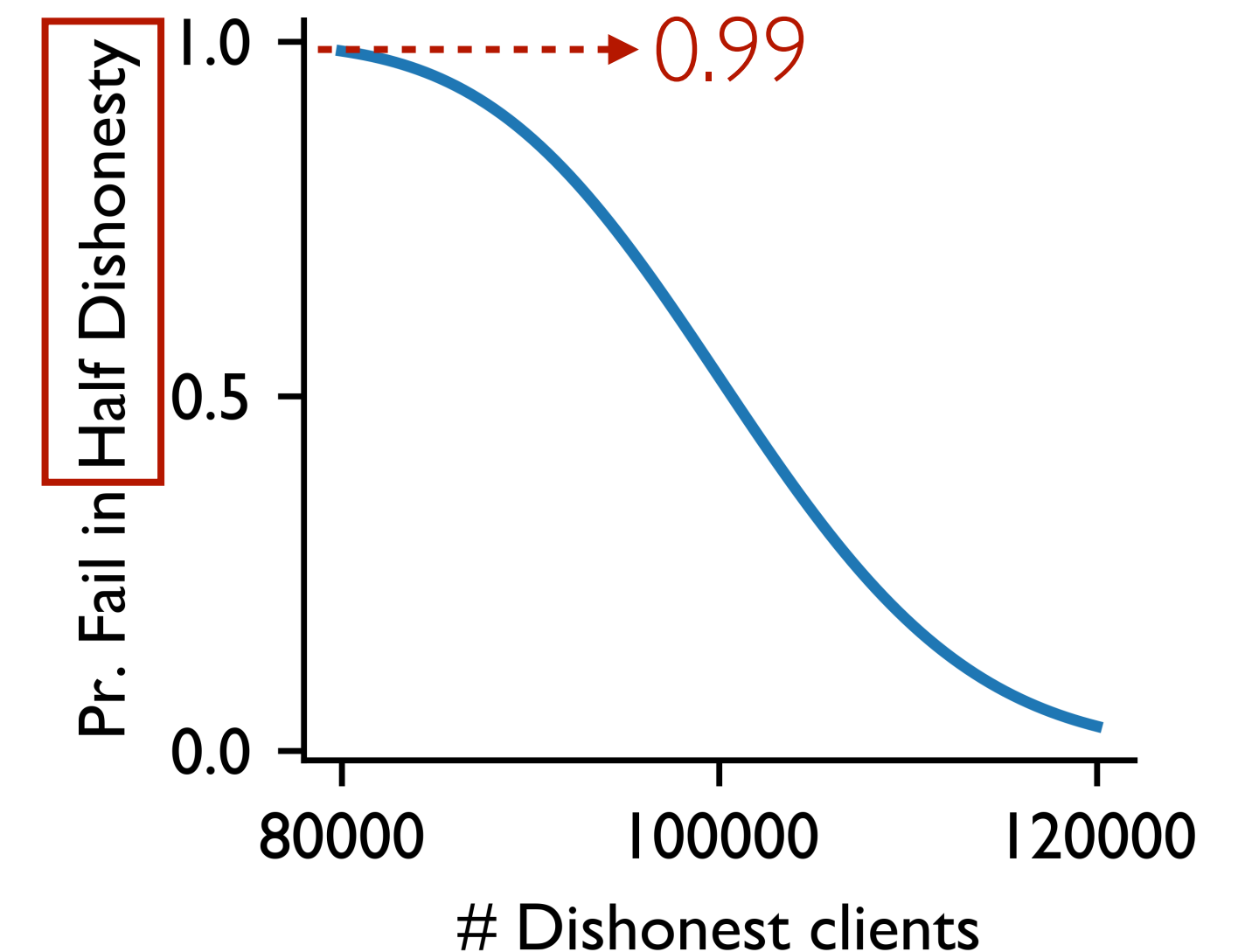
Each participant
sees a list of peers who
presents only by chance.

↘ The absent will not get
arbitrarily ignored

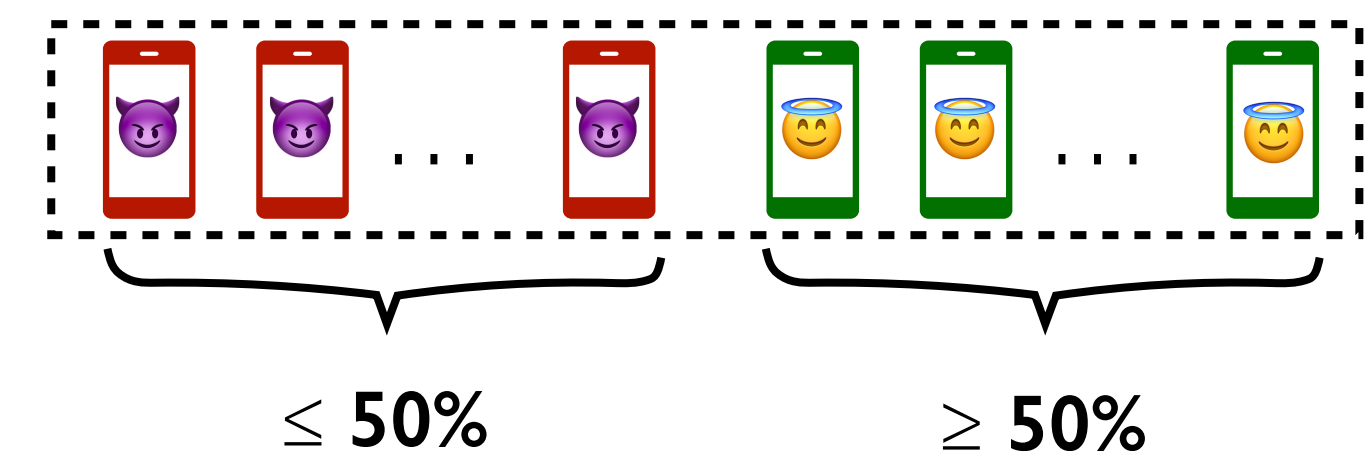
Solution: Enforce a **large enough list**
and a **small enough chance.**

Example

- **len(list):** ≥ 200
- **Chance:** $\leq 0.1\%$



Selected



Lotto: Random selection

What is achieved:

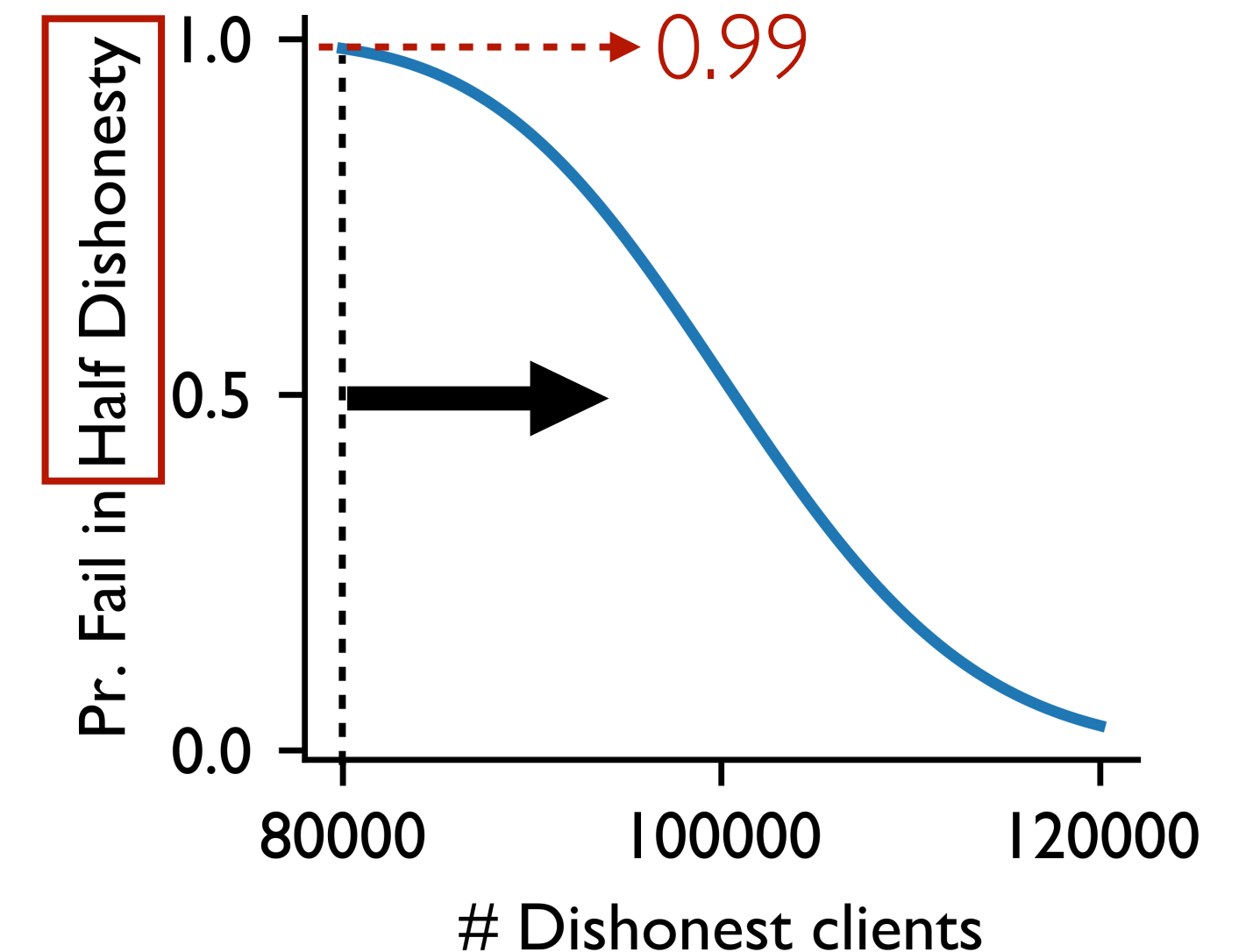
Each participant
sees a list of peers who
presents only by chance.

↘ The absent will not get
arbitrarily ignored

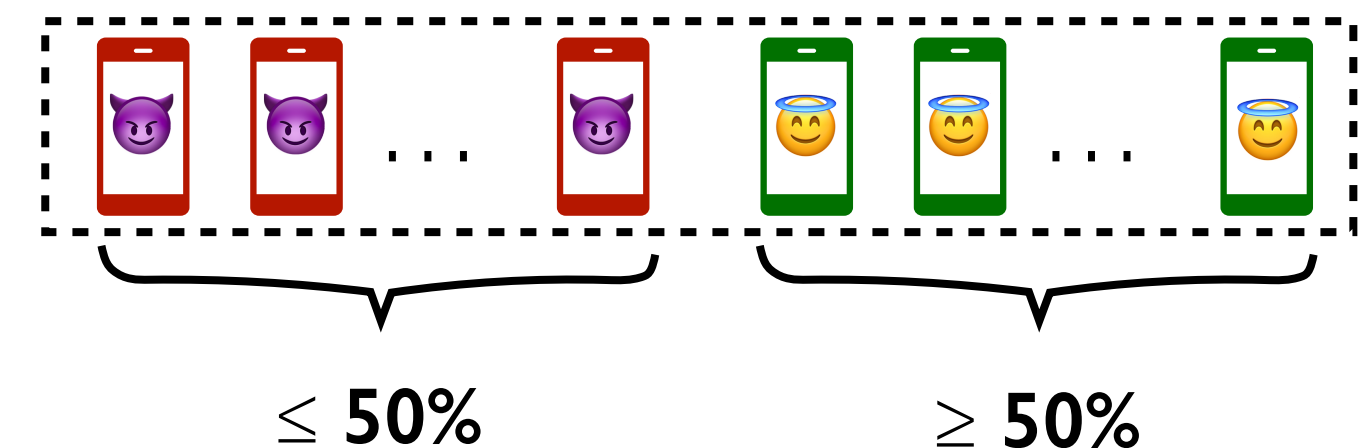
Solution: Enforce a **large enough list**
and a **small enough chance.**

Example

- **len(list):** ≥ 200
- **Chance:** $\leq 0.1\%$



Selected



Lotto: Random selection

What is achieved:

Each participant

sees a list of peers who
presents only by chance.

Predictable
to server?



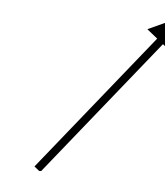
The absent will not get
arbitrarily ignored

Examples: #2 will be selected as $\mathbf{RF}_{pk_2}(2) = 1 < 3$.

Public Round index



Public Public keys



Lotto: Random selection

What is achieved:

Each participant

sees a list of peers who
presents only by chance.

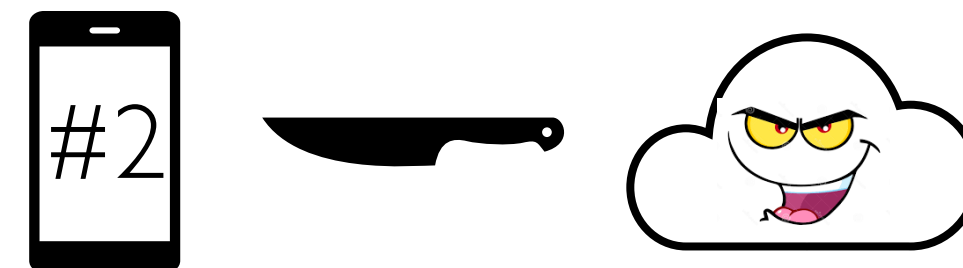
Predictable
to server?



The absent will not get
arbitrarily ignored

Problem: Attack surfaces **enlarged!**

Examples: #2 will be selected as $\mathbf{RF}_{pk2}(2) = 1 < 3$.
Before training, the server may grow its advantage by



Focused hacking

Lotto: Random selection

What is achieved:

Each participant

sees a list of peers who
presents only by chance.

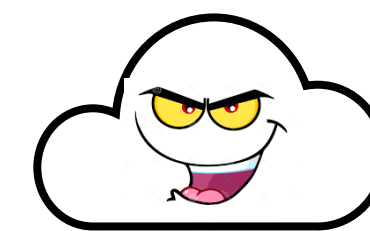
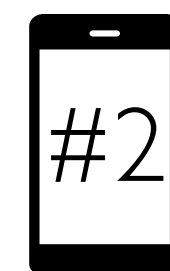
Predictable
to server?



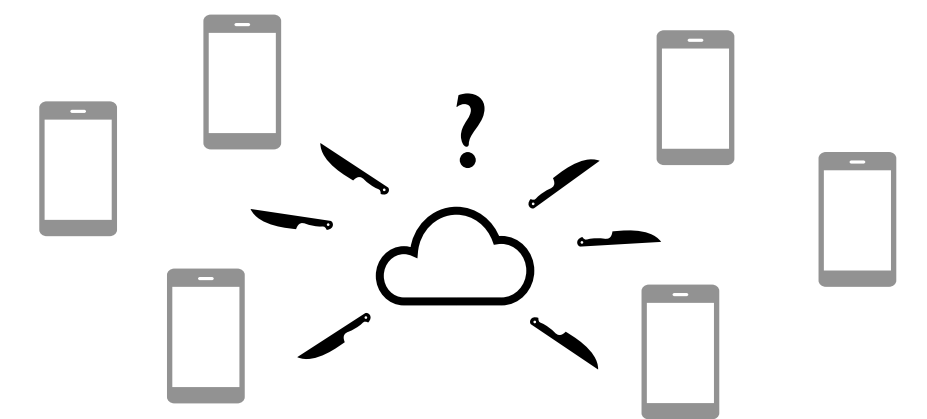
The absent will not get
arbitrarily ignored

Problem: Attack surfaces **enlarged!**

Examples: #2 will be selected as $\mathbf{RF}_{pk2}(2) = 1 < 3$.
Before training, the server may grow its advantage by



vs



Focused hacking

Random compromise

Lotto: Random selection

What is achieved:

Each participant

sees a list of peers who
presents only by chance.

Predictable
to server?

↘ The absent will not get
arbitrarily ignored

Solution: Self-sampling with
verifiable random functions (**VRFs**)^{1,2}.



Evaluation: $\mathbf{VRF.eval}_{sk2}(2) = (l, \dots)$ (output, ...)

Secret key ↗

¹Micali et al. "Verifiable random functions", In FOCS '99

²Dodis et al. "A verifiable random function with short proofs and keys", In PKC '05

Lotto: Random selection

What is achieved:

Each participant

sees a list of peers who
presents only by chance.

Predictable
to server?



The absent will not get
arbitrarily ignored

Solution: Self-sampling with
verifiable random functions (**VRFs**)^{1,2}.



Evaluation: $\mathbf{VRF.eval}_{sk_2}(2) = (1, \boldsymbol{\pi}_2)$ (output, **proof**)

¹Micali et al. "Verifiable random functions", In FOCS '99

²Dodis et al. "A verifiable random function with short proofs and keys", In PKC '05

Lotto: Random selection

What is achieved:

Each participant

sees a list of peers who
presents only by chance.

Predictable
to server?

↘ The absent will not get
arbitrarily ignored

Solution: Self-sampling with
verifiable random functions (**VRFs**)^{1,2}.



Evaluation: $\mathbf{VRF.eval}_{sk_2}(2) = (l, \boldsymbol{\pi}_2)$ (output, **proof**)

Verification: $\mathbf{VRF.ver}_{pk_2}(2, l, \boldsymbol{\pi}_2) = \text{True}$

Public key ↗

¹Micali et al. "Verifiable random functions", In FOCS '99

²Dodis et al. "A verifiable random function with short proofs and keys", In PKC '05

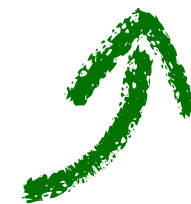
Lotto: Random selection

What is achieved:

Each participant

sees a list of peers who
presents only by chance.

Unpredictable
to server



The absent will not get
arbitrarily ignored

Solution: Self-sampling with
verifiable random functions (**VRFs**)^{1,2}.

I self-sample
with (I, π_2)



Evaluation: $\mathbf{VRF.eval}_{sk_2}(2) = (I, \pi_2)$ (output, **proof**)

Verification: $\mathbf{VRF.ver}_{pk_2}(2, I, \pi_2) = \text{True}$

¹Micali et al. "Verifiable random functions", In FOCS '99

²Dodis et al. "A verifiable random function with short proofs and keys", In PKC '05

Lotto: Random selection

What is achieved:

Each participant

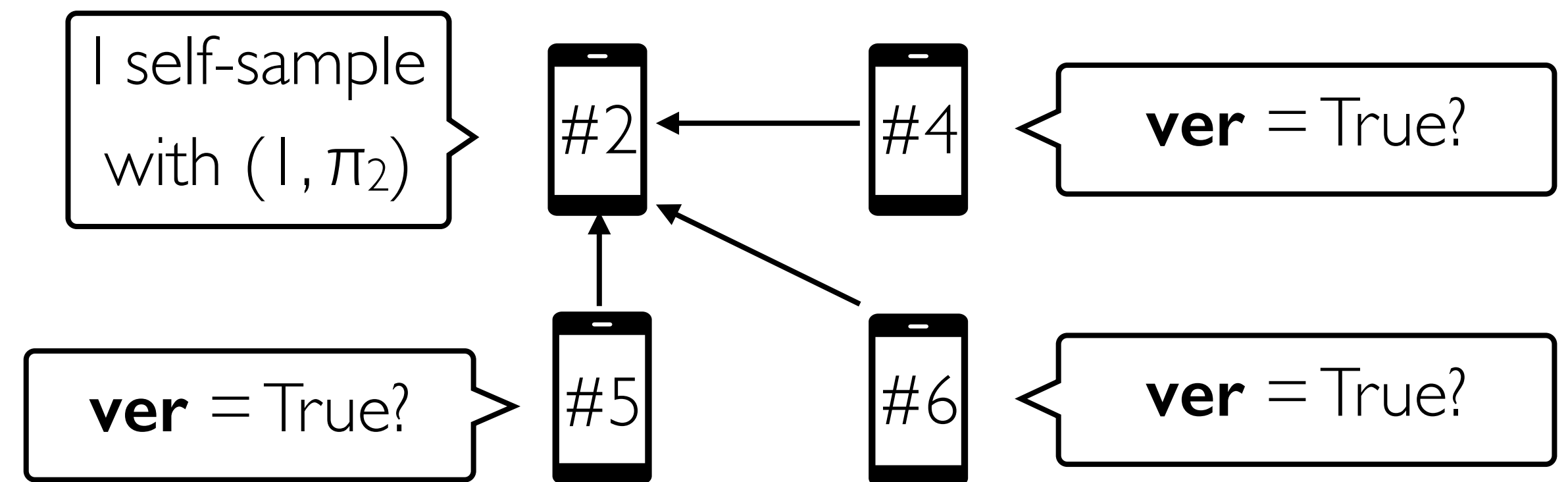
sees a list of peers who
presents only by chance.

Unpredictable
to server



The absent will not get
arbitrarily ignored

Solution: Self-sampling with
verifiable random functions (**VRFs**)^{1,2}.



Evaluation: $\mathbf{VRF.eval}_{sk_2}(2) = (l, \boldsymbol{\pi}_2)$ (output, **proof**)

Verification: $\mathbf{VRF.ver}_{pk_2}(2, l, \boldsymbol{\pi}_2) = \text{True}$

¹Micali et al. "Verifiable random functions", In FOCS '99

²Dodis et al. "A verifiable random function with short proofs and keys", In PKC '05

Lotto: Random selection

What is achieved:

Each participant

sees a list of peers who
presents only by chance.

Unpredictable
to server



The absent will not get
arbitrarily ignored

Minor issues:

- **Participant consistency:** leverage SecAgg
- **Fixed sample size:** over-selection
- **Consistent round index:** uniqueness check

...

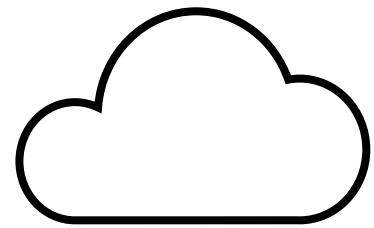
Please find more in the paper :)

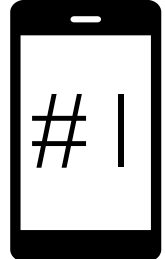
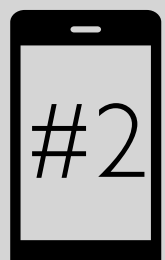
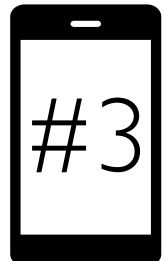
¹Thus also of distributed DP (other privacy-enhancing techniques may not have this feature and this is left for future work).

Lotto: Informed selection

Lotto: Informed selection

Example

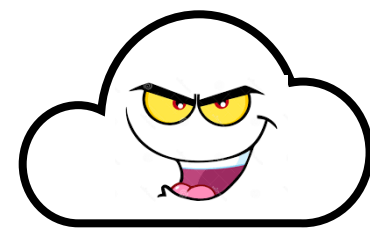
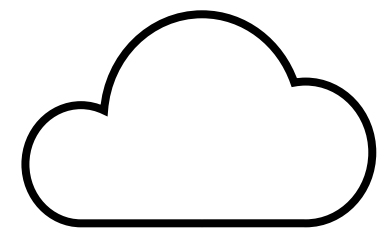


	(Est.) latency	Select	(Est.) latency	Select
	1.2s	Yes		Yes
	2.7s	No	Does NOT matter.	No
	1.6s	Yes		No
...

Selection criteria: the fastest For dishonest majority

Lotto: Informed selection

Example



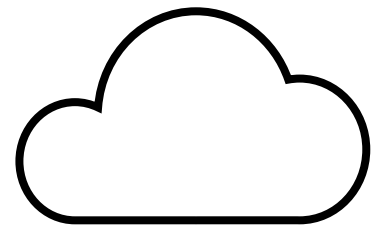
	(Est.) latency	Select	(Est.) latency	Select
#1	1.2s	Yes		Yes
#2	2.7s	No	Does NOT matter.	No
#3	1.6s	Yes		No
...

Selection criteria: the fastest For dishonest majority

Major Challenge: Client metrics are **hard to verify** by honest clients

Lotto: Informed selection

Example

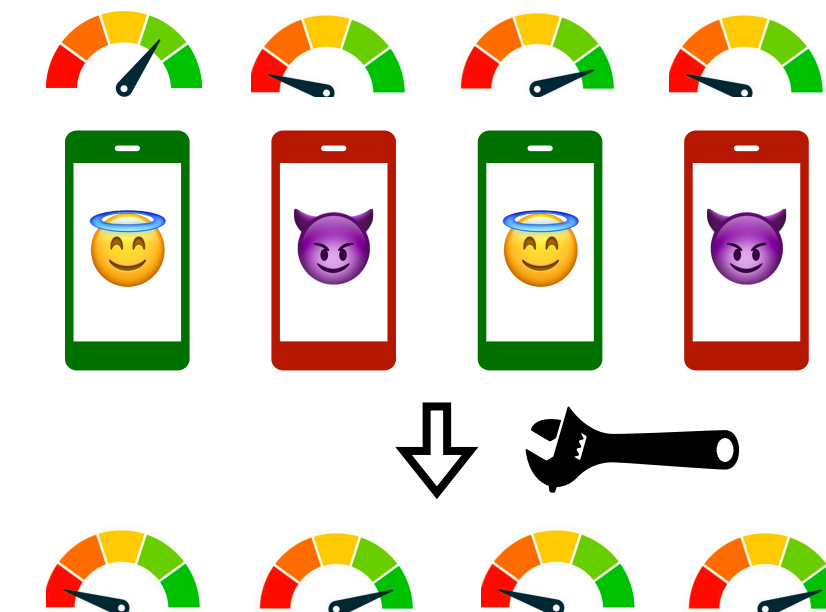


	(Est.) latency	Select	(Est.) latency	Select
#1	1.2s	Yes		Yes
#2	2.7s	No	Does NOT matter.	No
#3	1.6s	Yes		No
...

Selection criteria: the fastest For dishonest majority

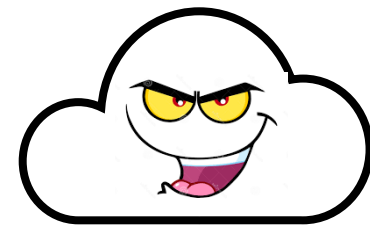
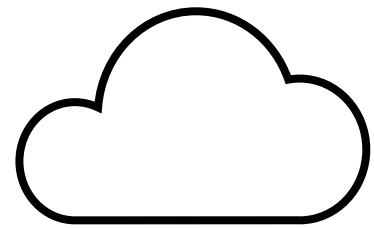
Major Challenge: Client metrics are **hard to verify** by honest clients

Metrics can be easily fake



Lotto: Informed selection

Example

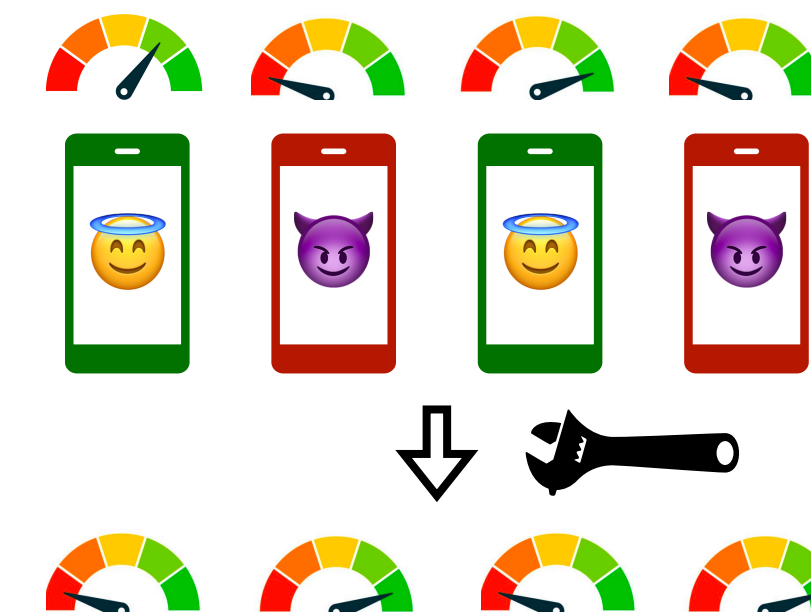


	(Est.) latency	Select	(Est.) latency	Select
#1	1.2s	Yes		Yes
#2	2.7s	No	Does NOT matter.	No
#3	1.6s	Yes		No
...

Selection criteria: the fastest For dishonest majority

Major Challenge: Client metrics are **hard to verify** by honest clients

Metrics can be easily fake



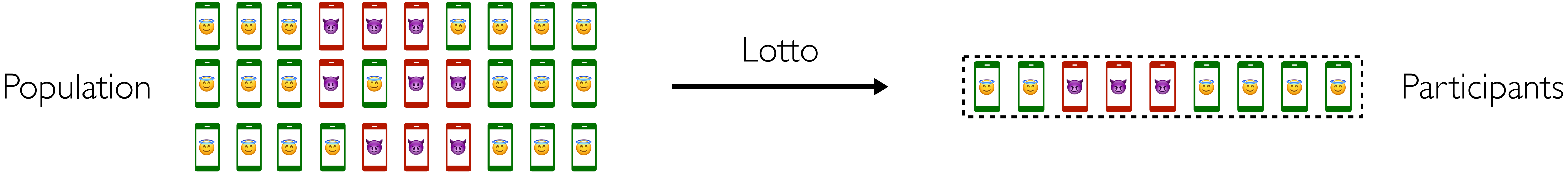
Solution: **Approximate** inform selection by **random** selection

Please find more in the paper :)

Lotto prevents arbitrary manipulation

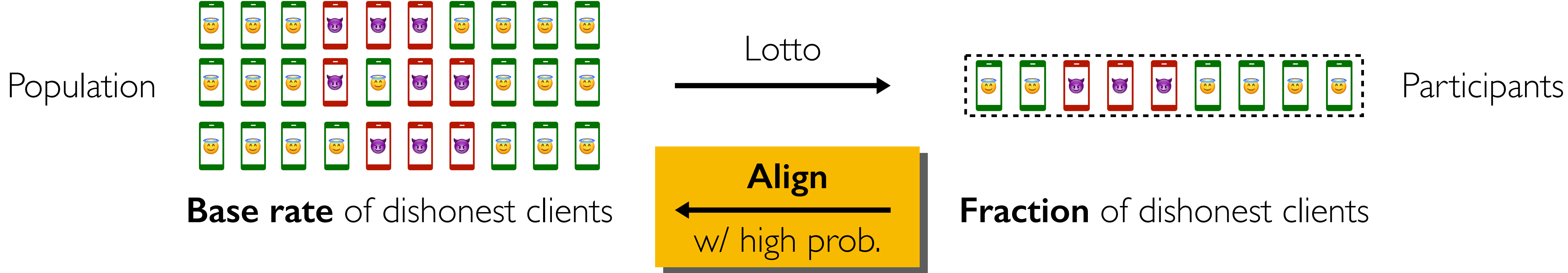
Lotto prevents arbitrary manipulation

What can be **proven**:



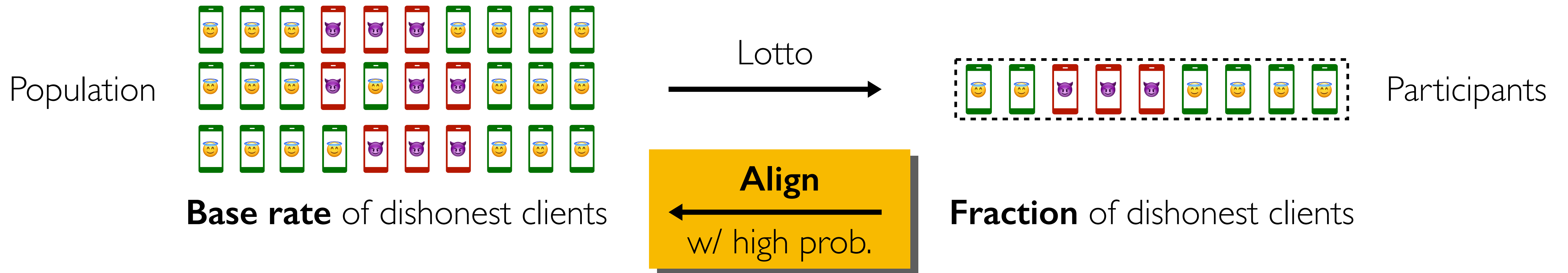
Lotto prevents arbitrary manipulation

What can be **proven**:



Lotto prevents arbitrary manipulation

What can be **proven**:

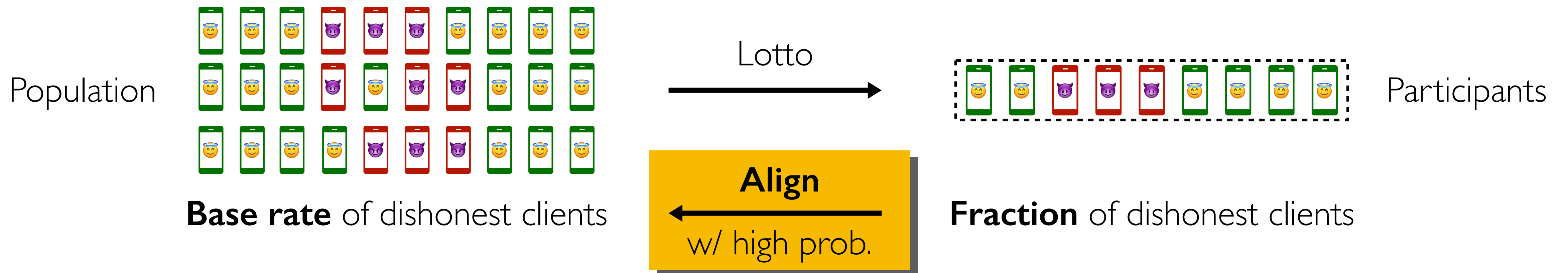


Example

- **Population:** 200,000
- **Dishonesty base rate:** 0.005

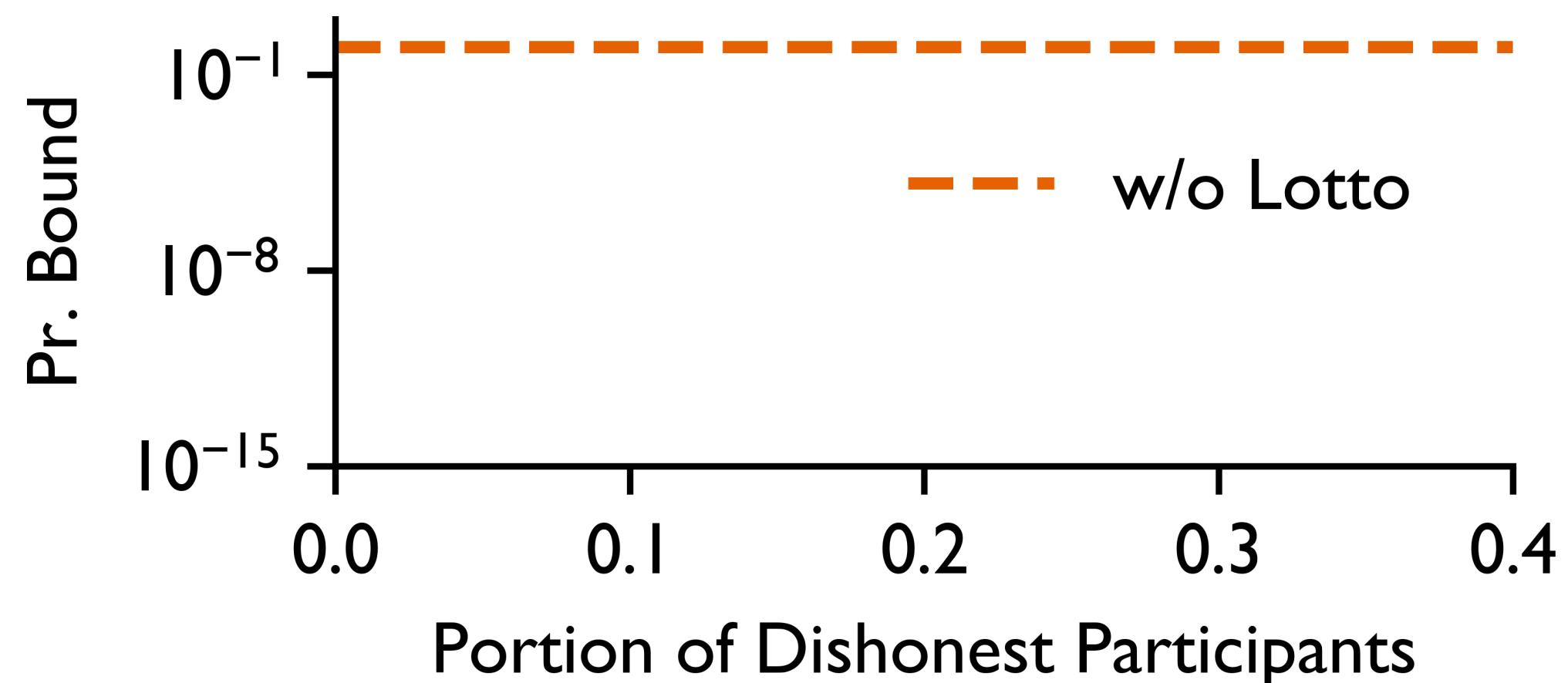
Lotto prevents arbitrary manipulation

What can be **proven**:



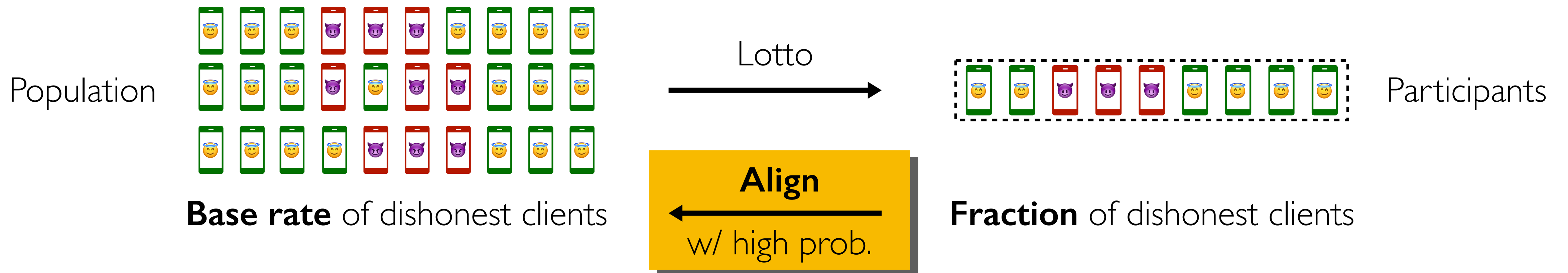
Example

- **Population:** 200,000
- **Dishonesty base rate:** 0.005
- **Target participants:** 200



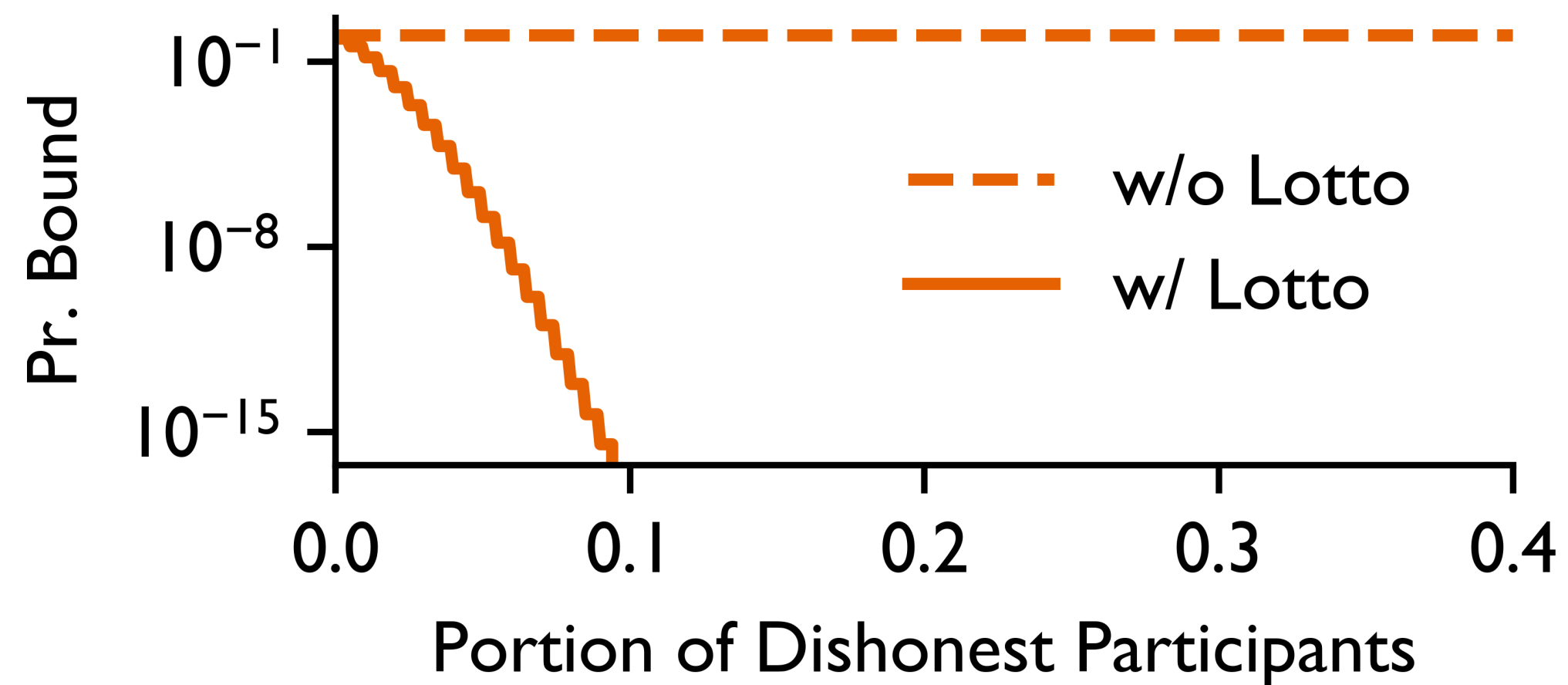
Lotto prevents arbitrary manipulation

What can be **proven**:



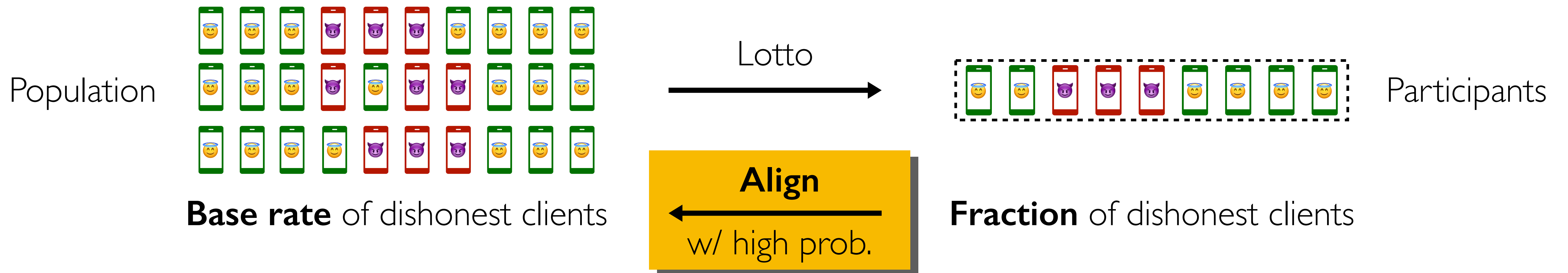
Example

- **Population:** 200,000
- **Dishonesty base rate:** 0.005
- **Target participants:** 200



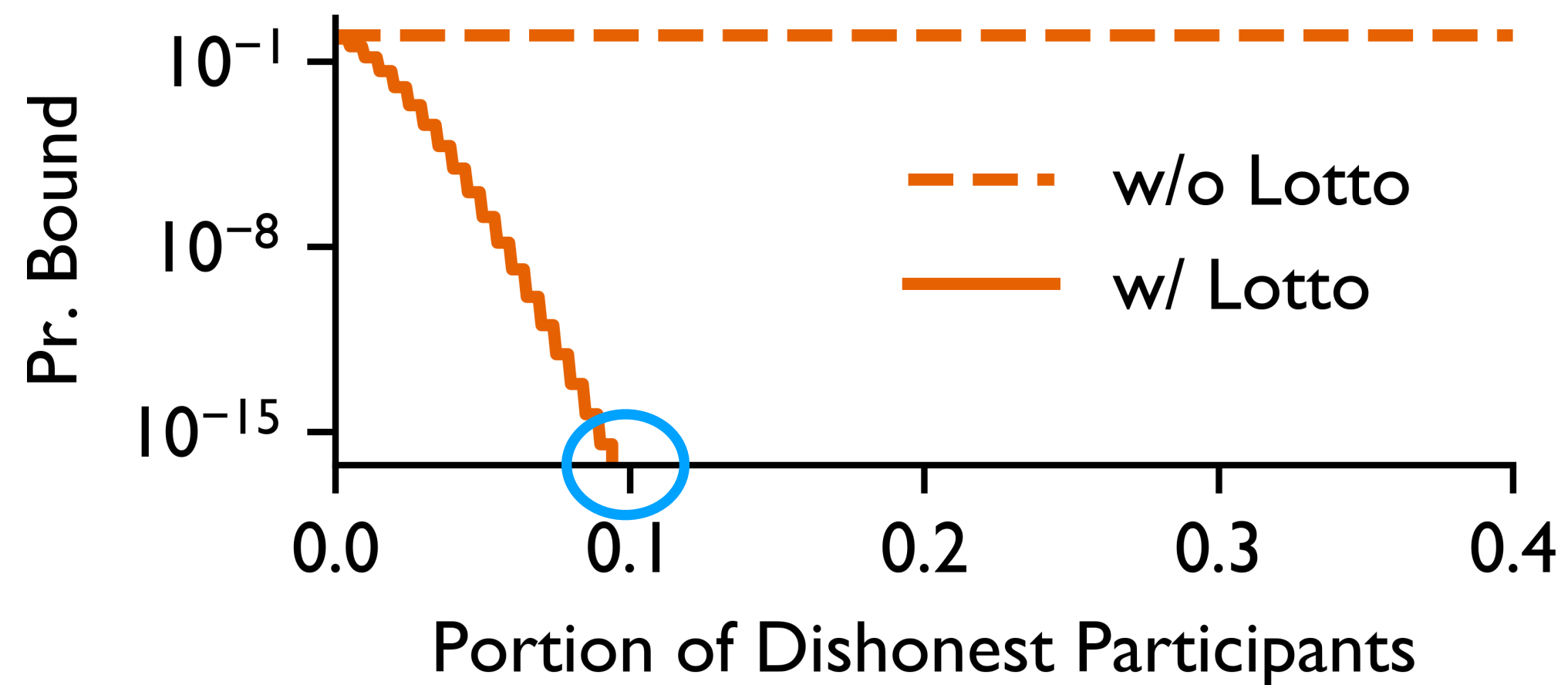
Lotto prevents arbitrary manipulation

What can be **proven**:



Example

- **Population:** 200,000
- **Dishonesty base rate:** 0.005
- **Target participants:** 200



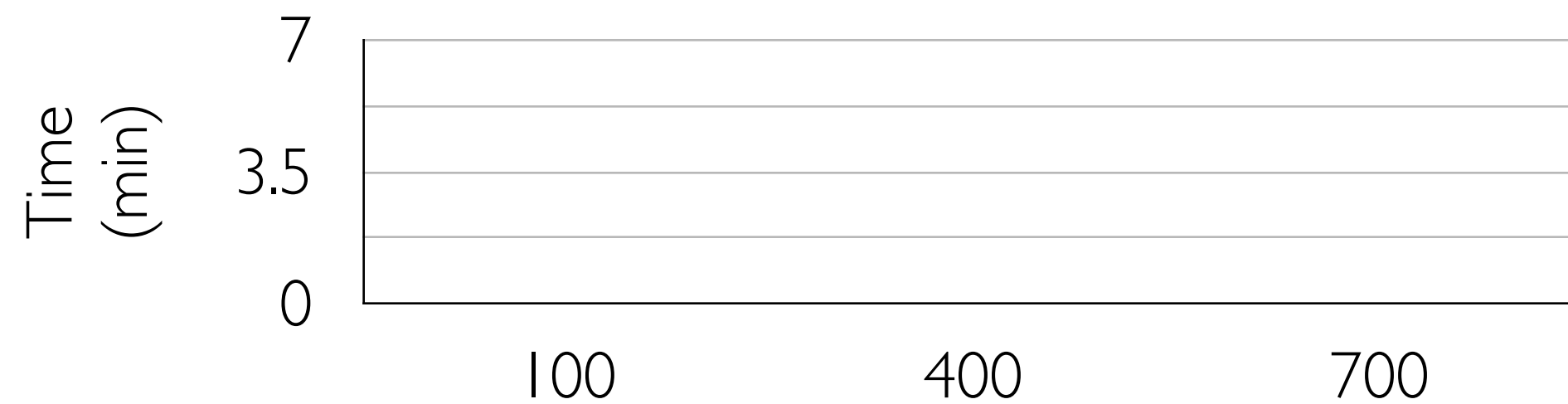
Lotto induces no or mild overhead

Lotto induces no or mild overhead

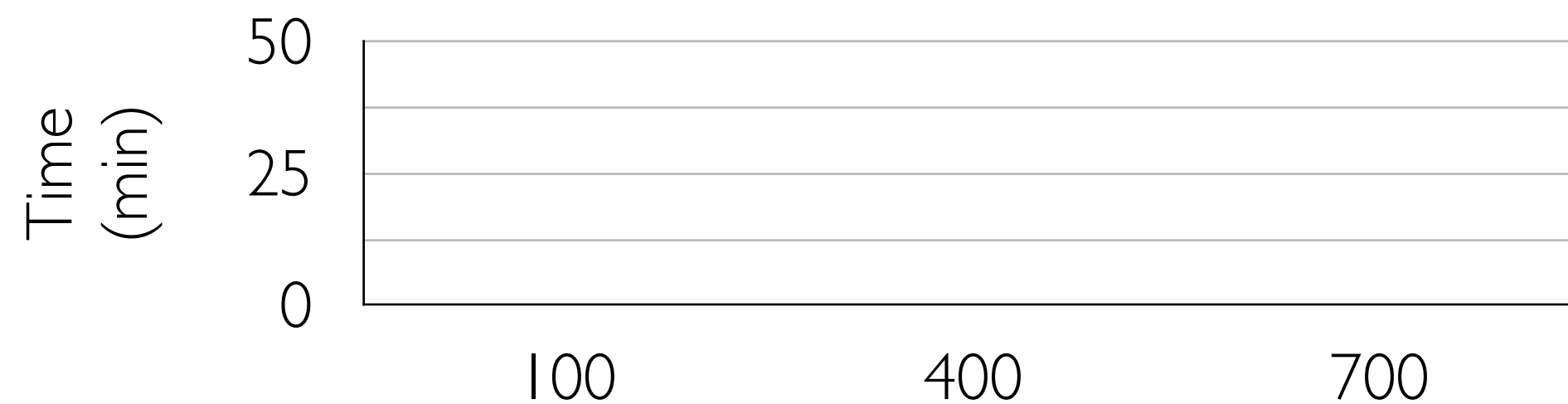
**FEMNIST
@CNN**



**OpenImage
@MobileNet**



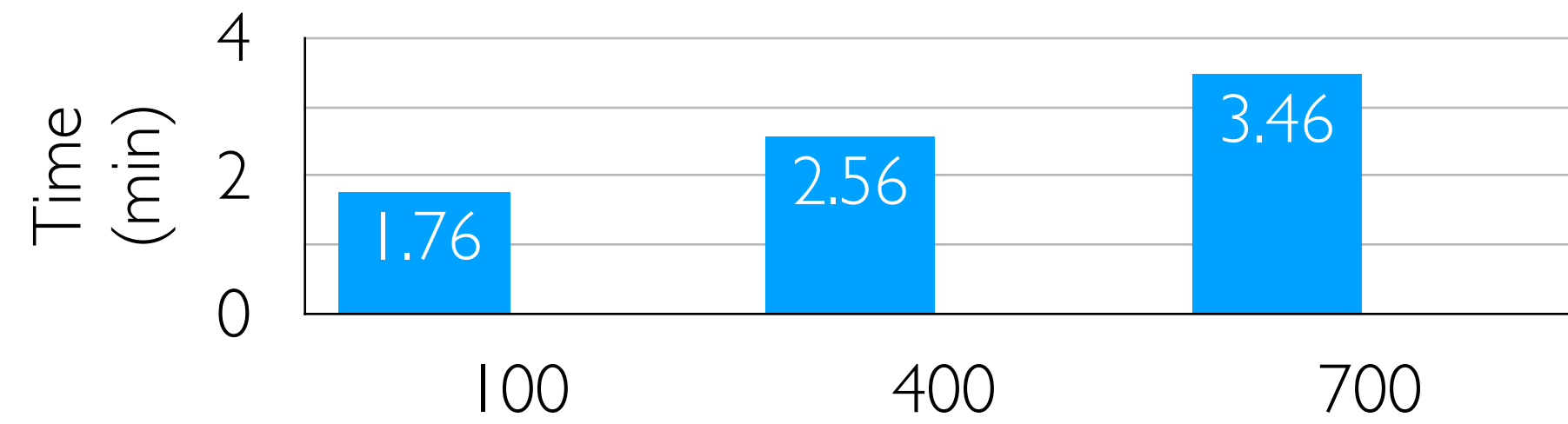
**Reddit
@Albert**



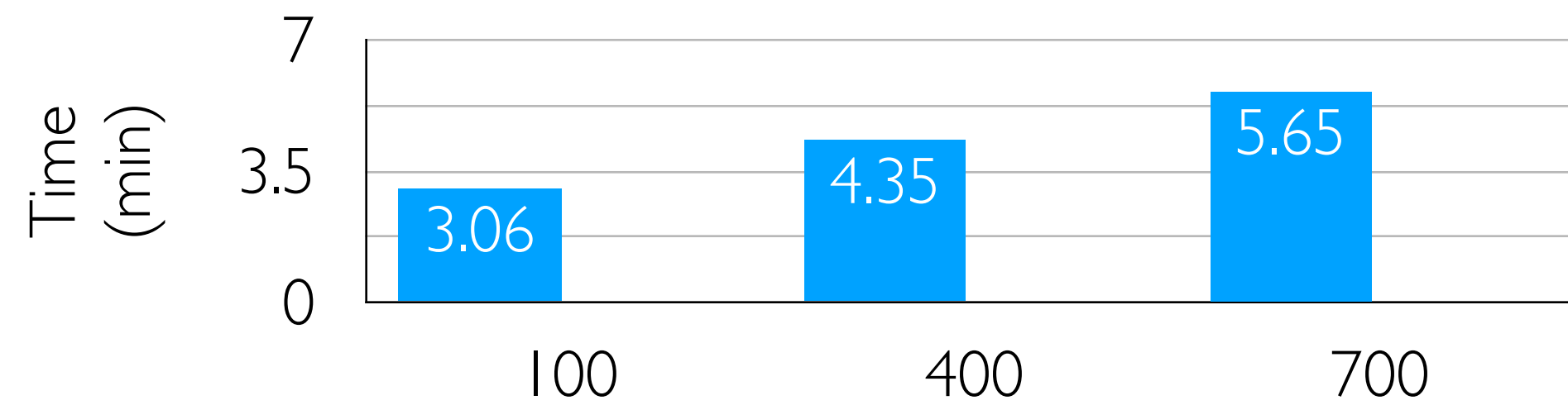
¹Random selection as an example. See results for informed selection in the paper.

Lotto induces no or mild overhead

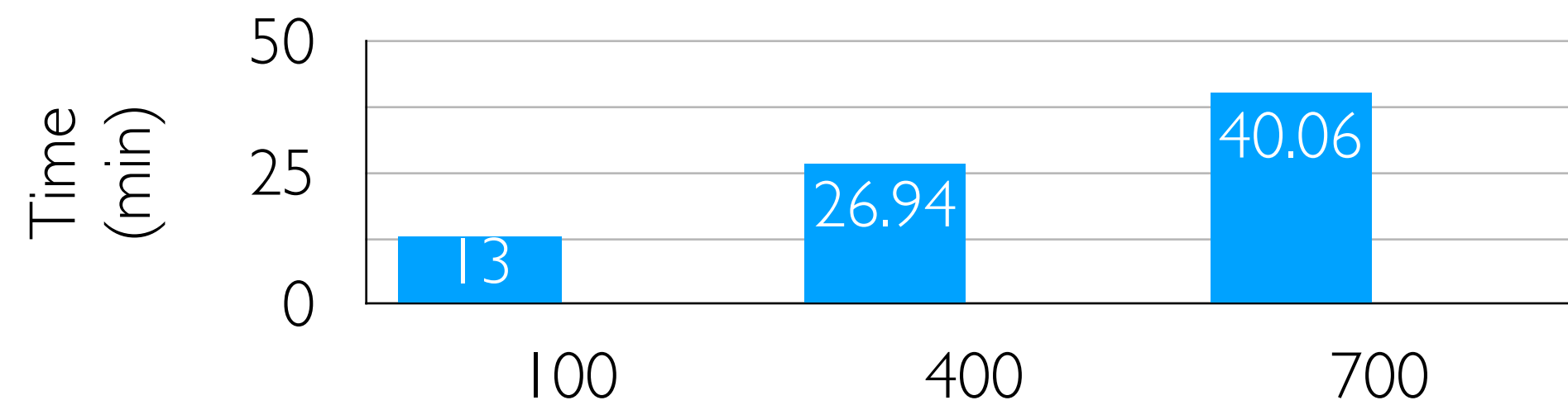
**FEMNIST
@CNN**



**OpenImage
@MobileNet**



**Reddit
@Albert**

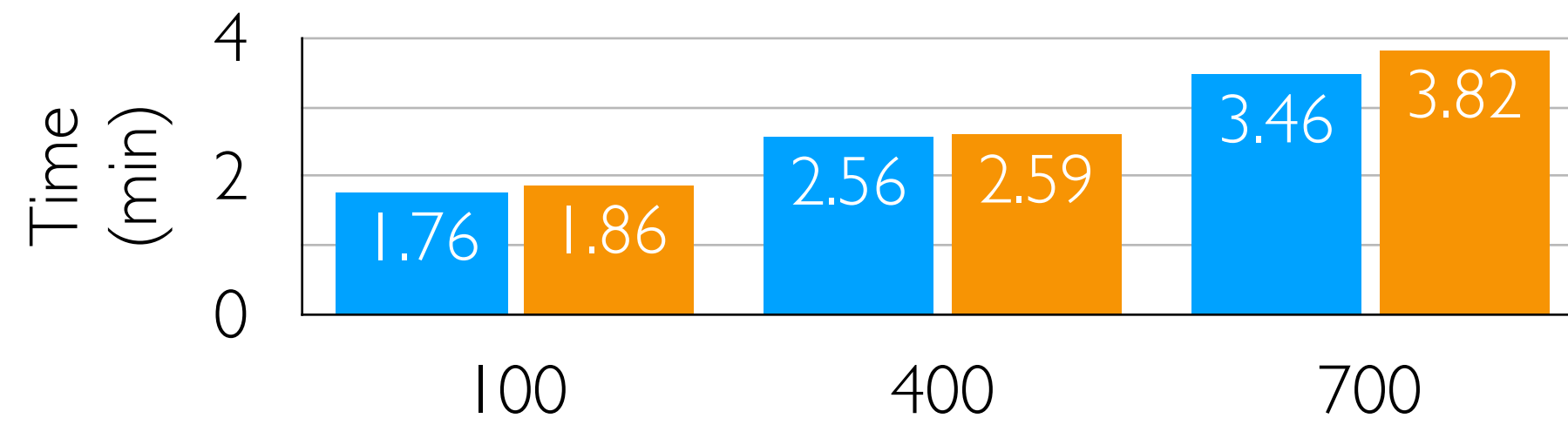


Population size

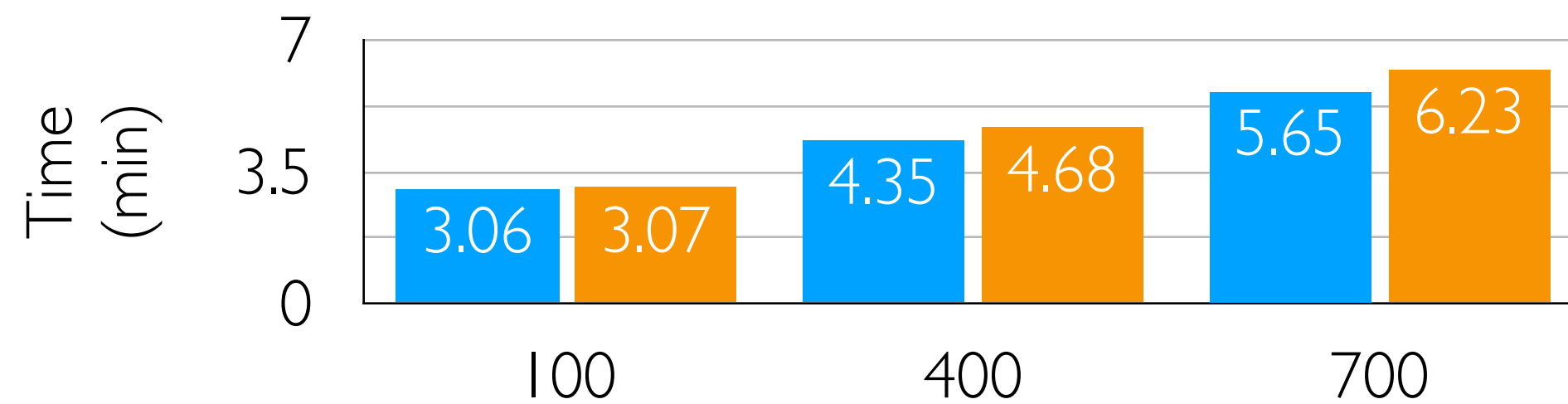
 **w/o Lotto**

Lotto induces no or mild overhead

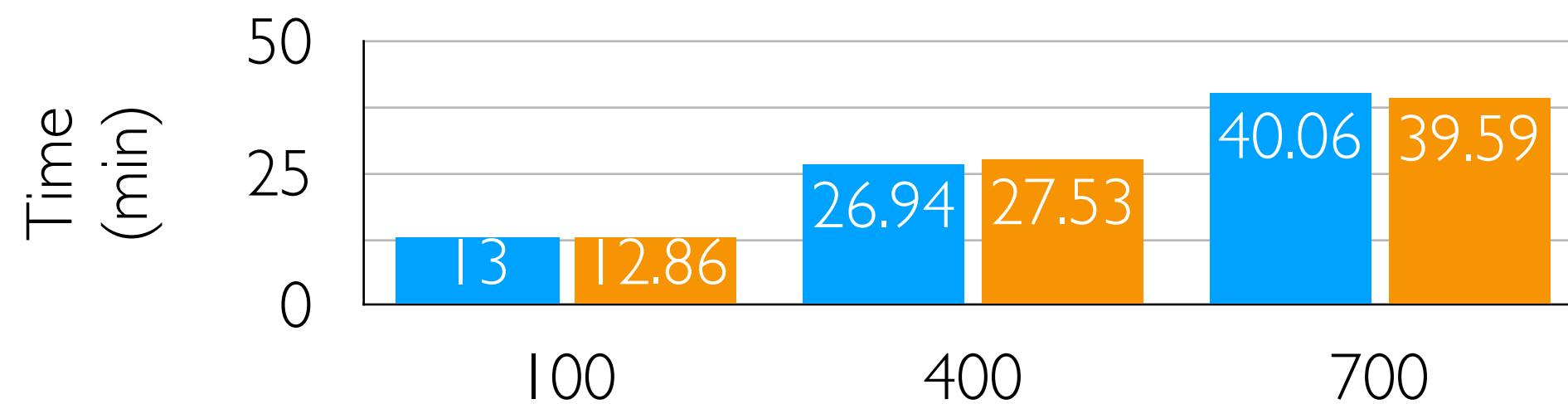
FEMNIST
@CNN



OpenImage
@MobileNet



Reddit
@Albert



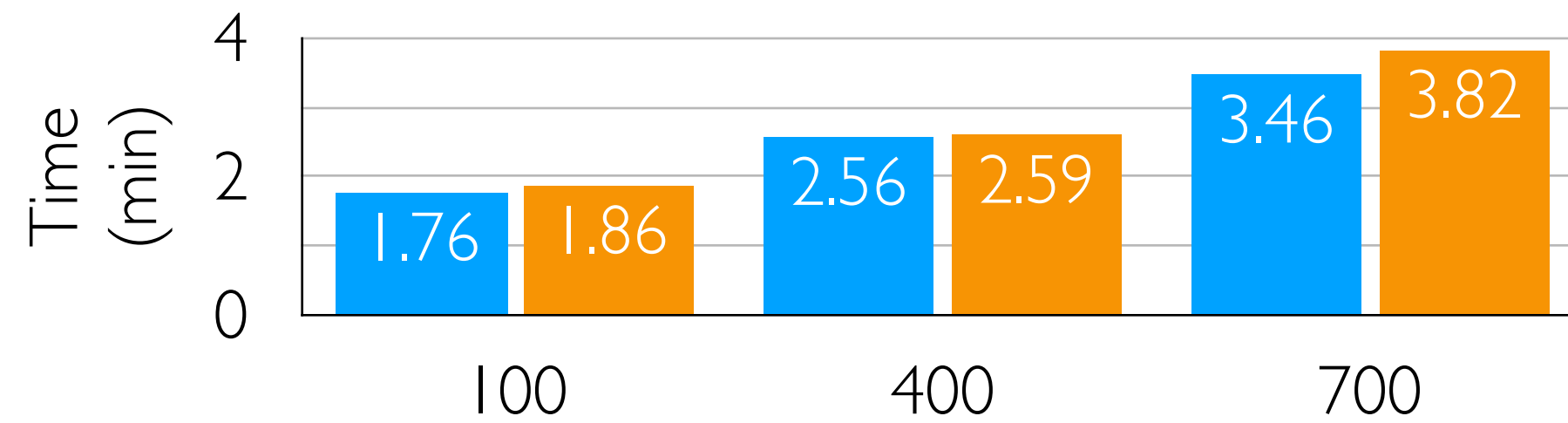
Population size

■ w/o Lotto
■ w/ Lotto

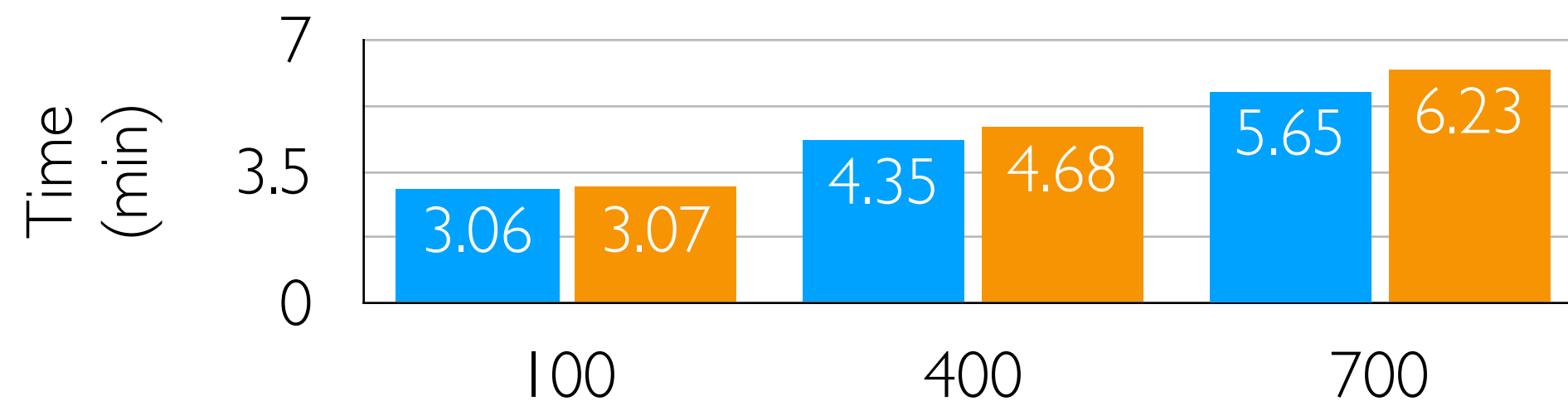
Lotto adds no more than **10%** in **time**

Lotto induces no or mild overhead

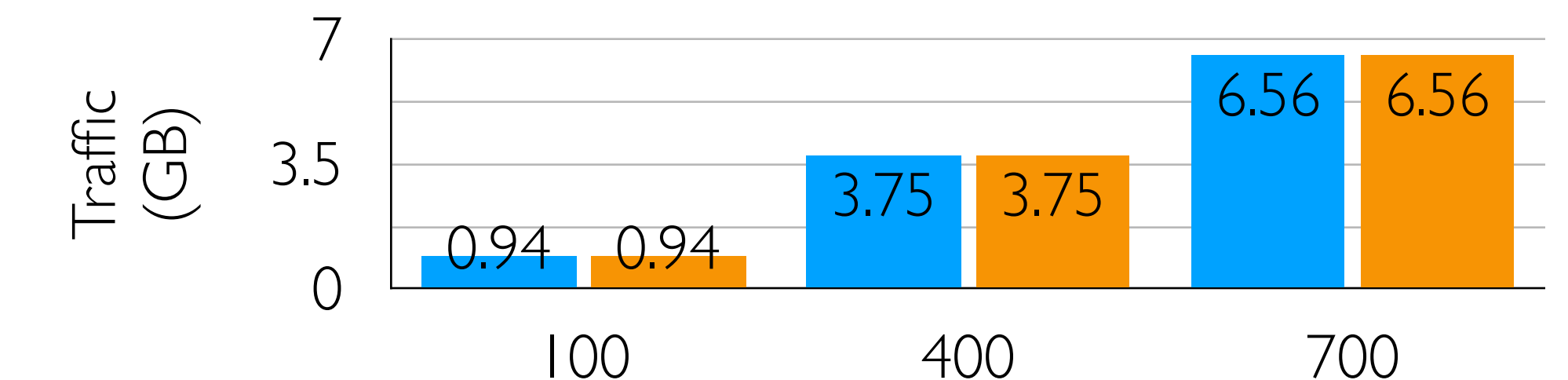
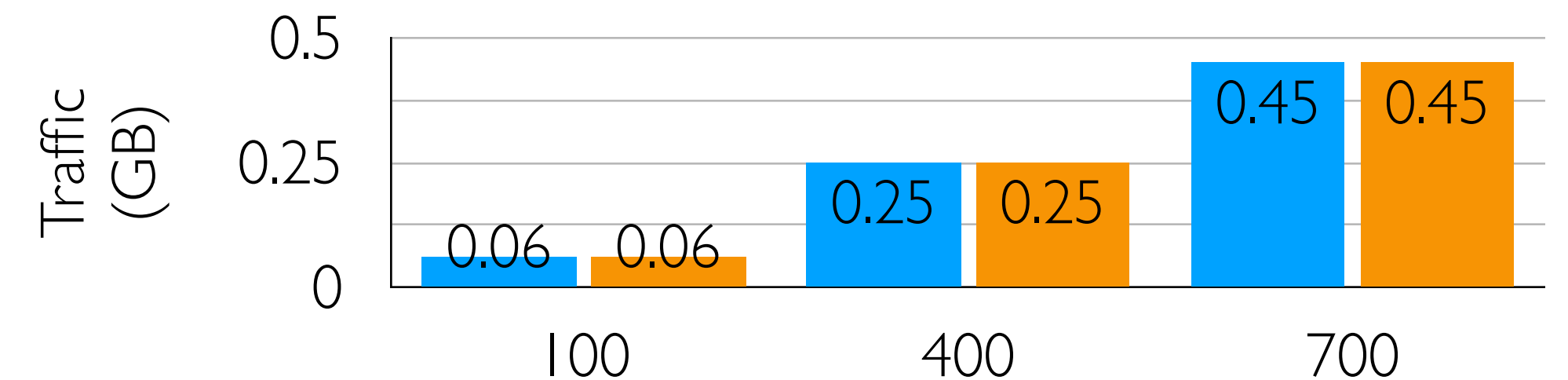
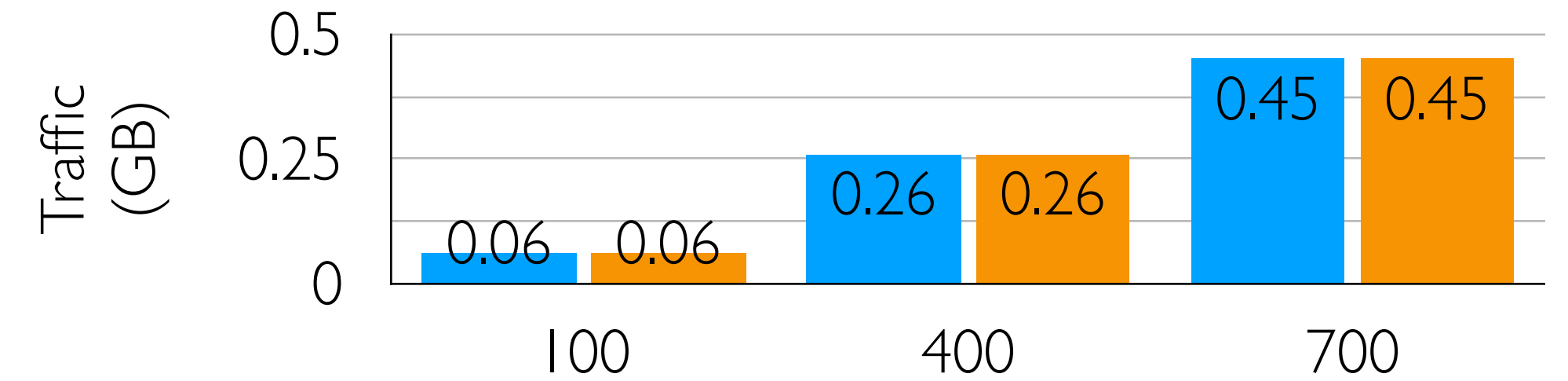
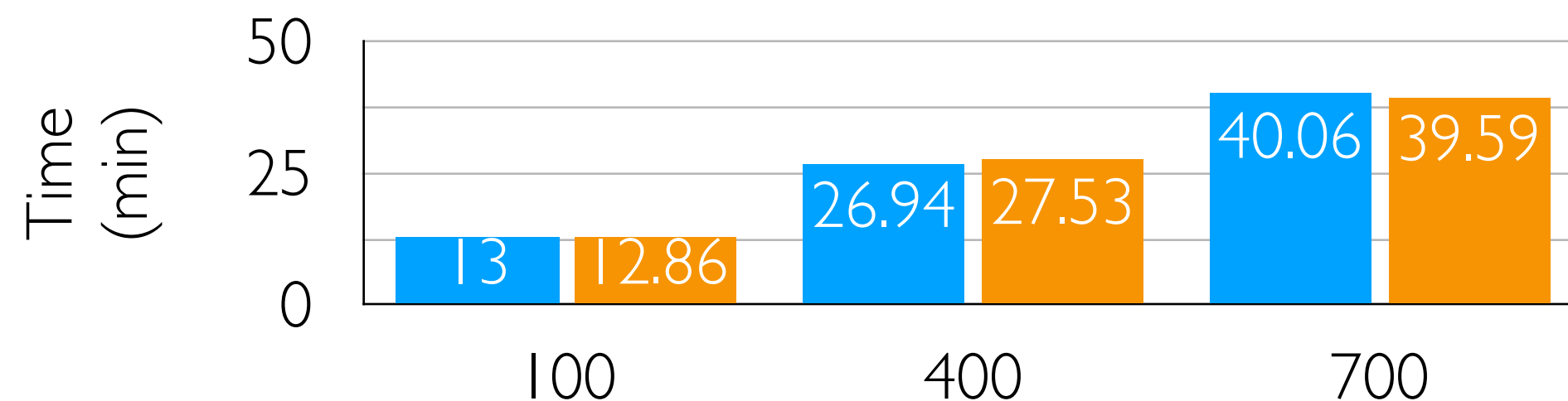
FEMNIST
@CNN



OpenImage
@MobileNet



Reddit
@Albert



■ w/o Lotto
■ w/ Lotto

Lotto adds no more than 10% in time

Lotto costs negligible in network

¹Random selection as an example. See results for informed selection in the paper.

Lotto functions as insecure selectors

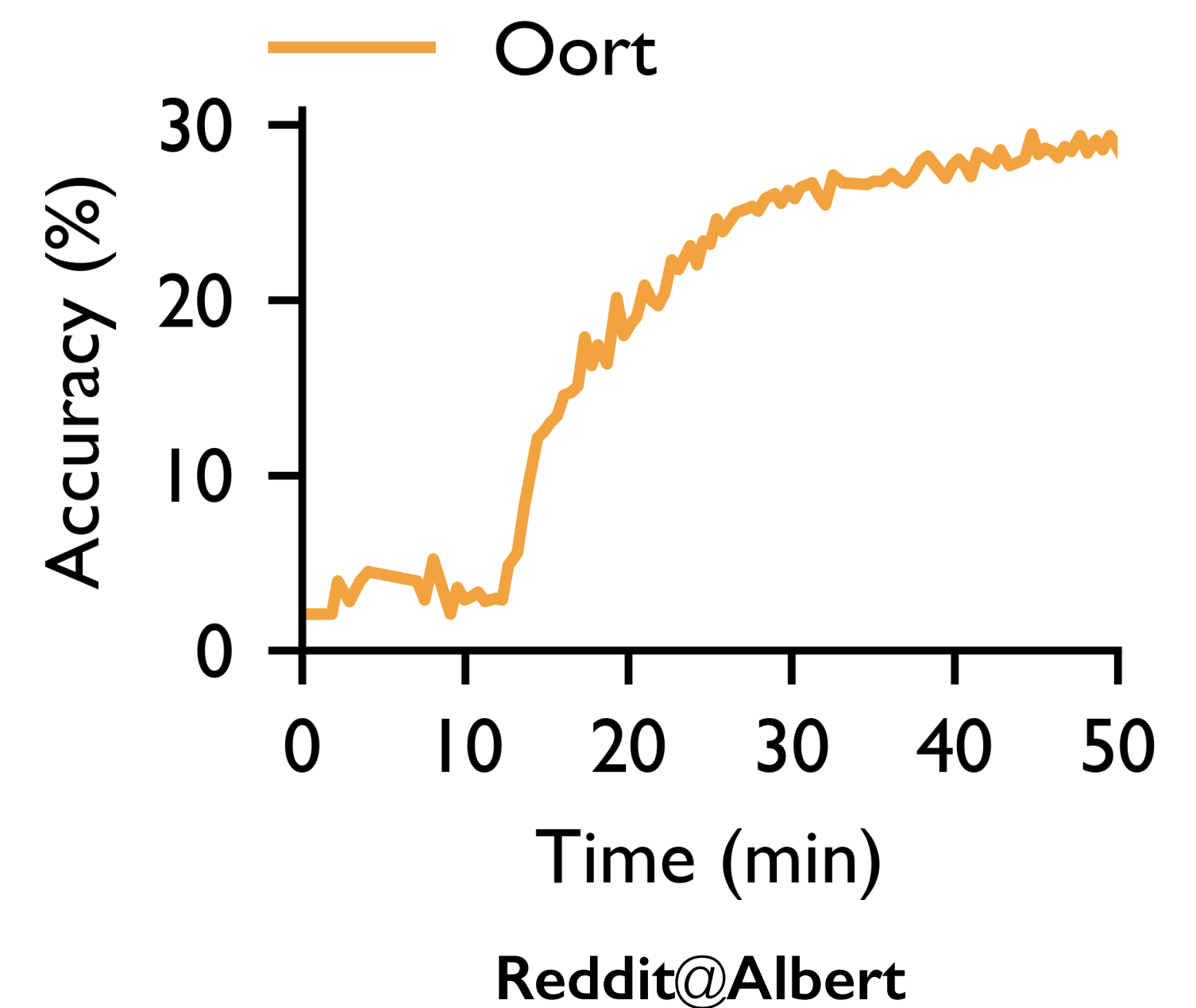
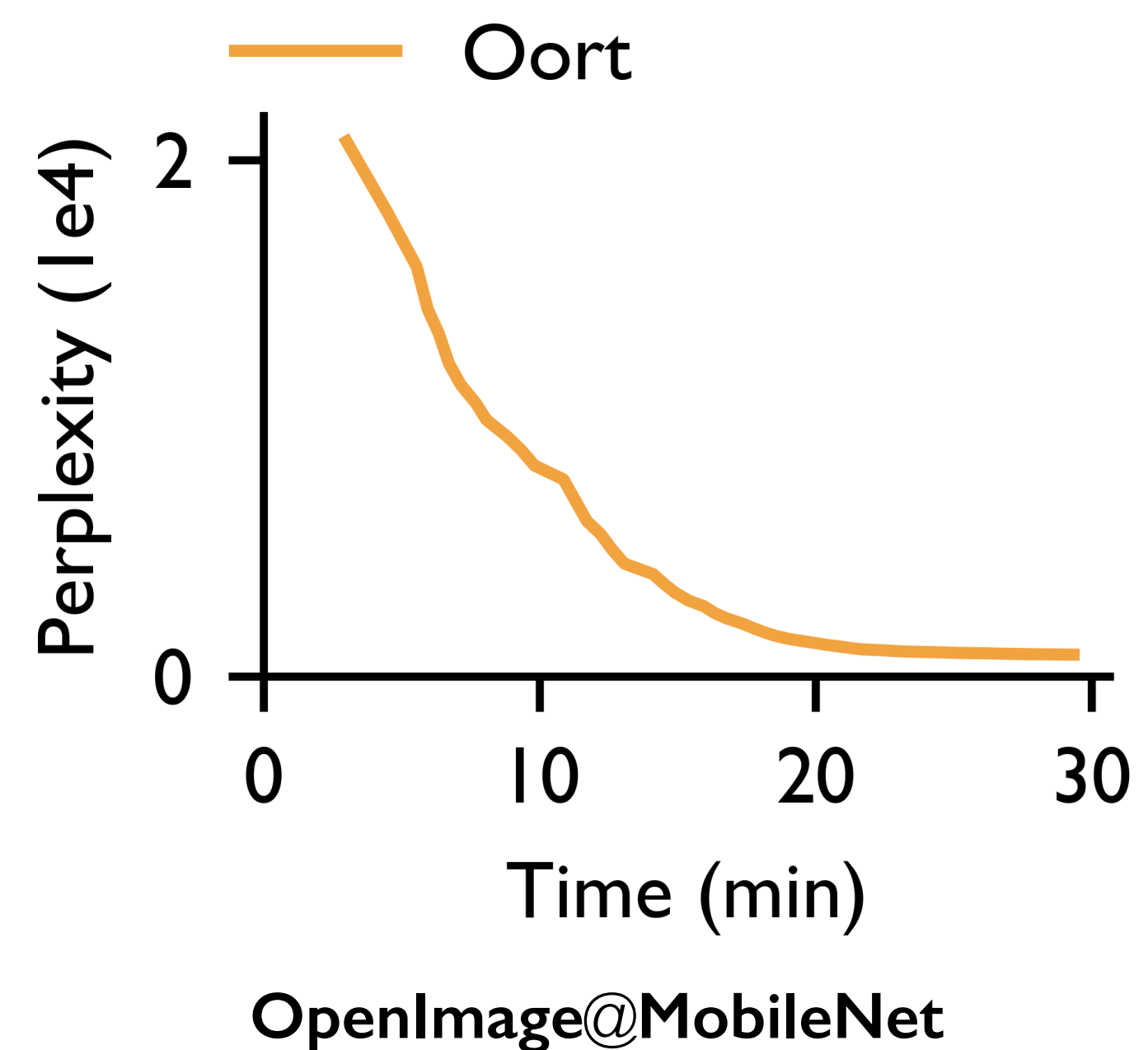
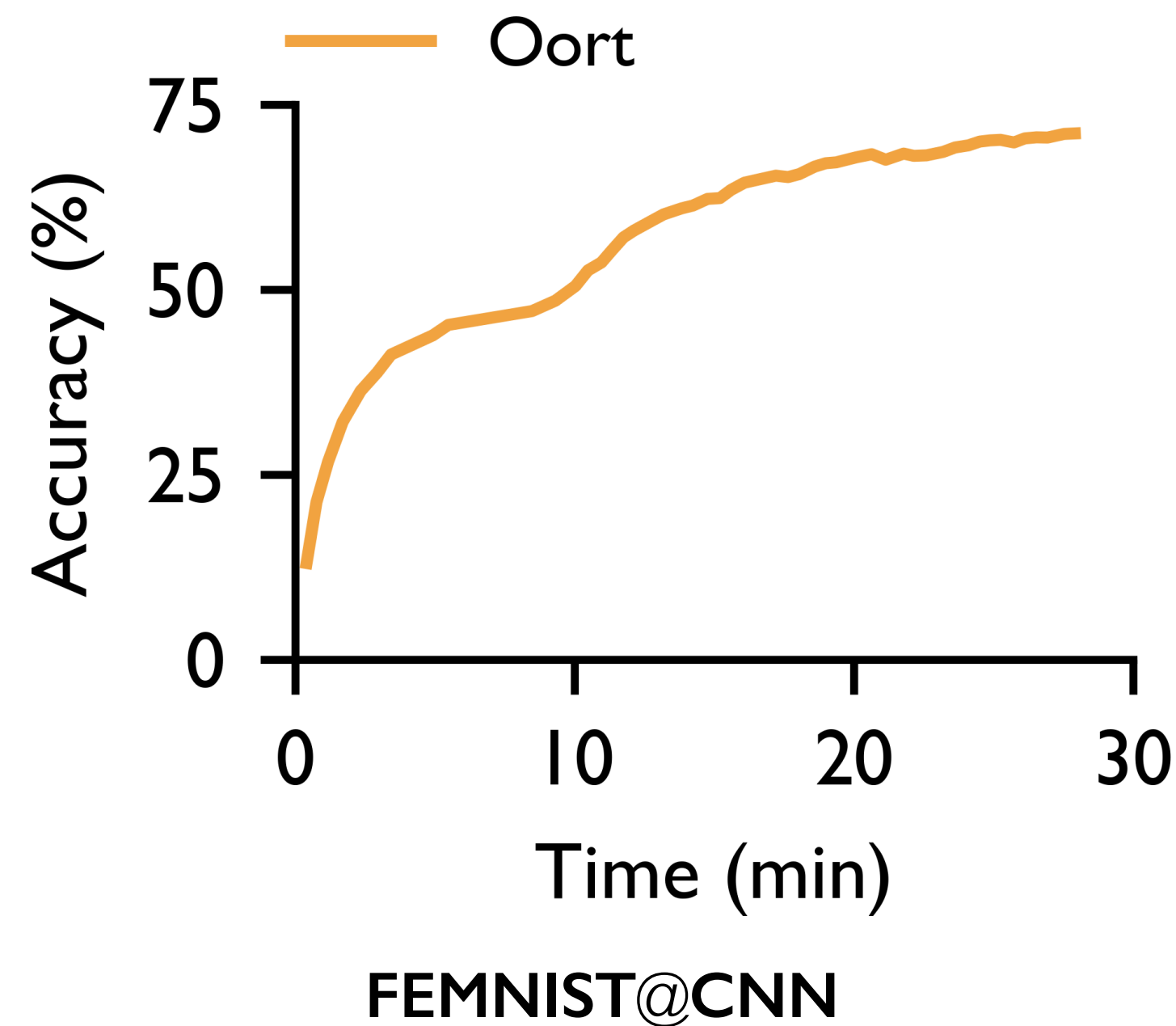
Lotto functions as insecure selectors

Oort¹ → State-of-the-art **informed** selector: optimized for **time-to-accuracy** of training

¹Lai et al. "Oort: Efficient Federated Learning via Guided Participant Selection", In OSDI '21

Lotto functions as insecure selectors

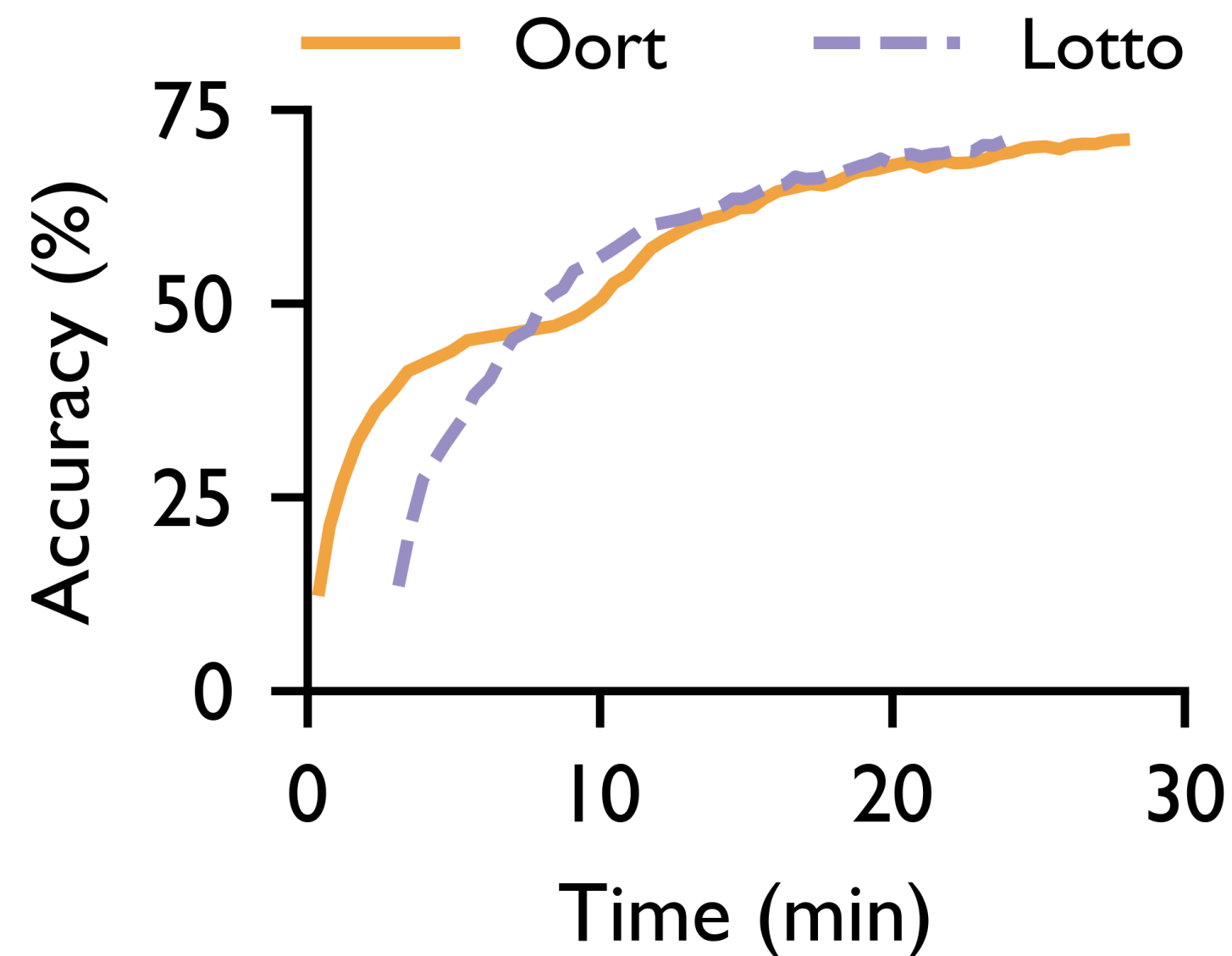
Oort¹ → State-of-the-art **informed** selector: optimized for **time-to-accuracy** of training



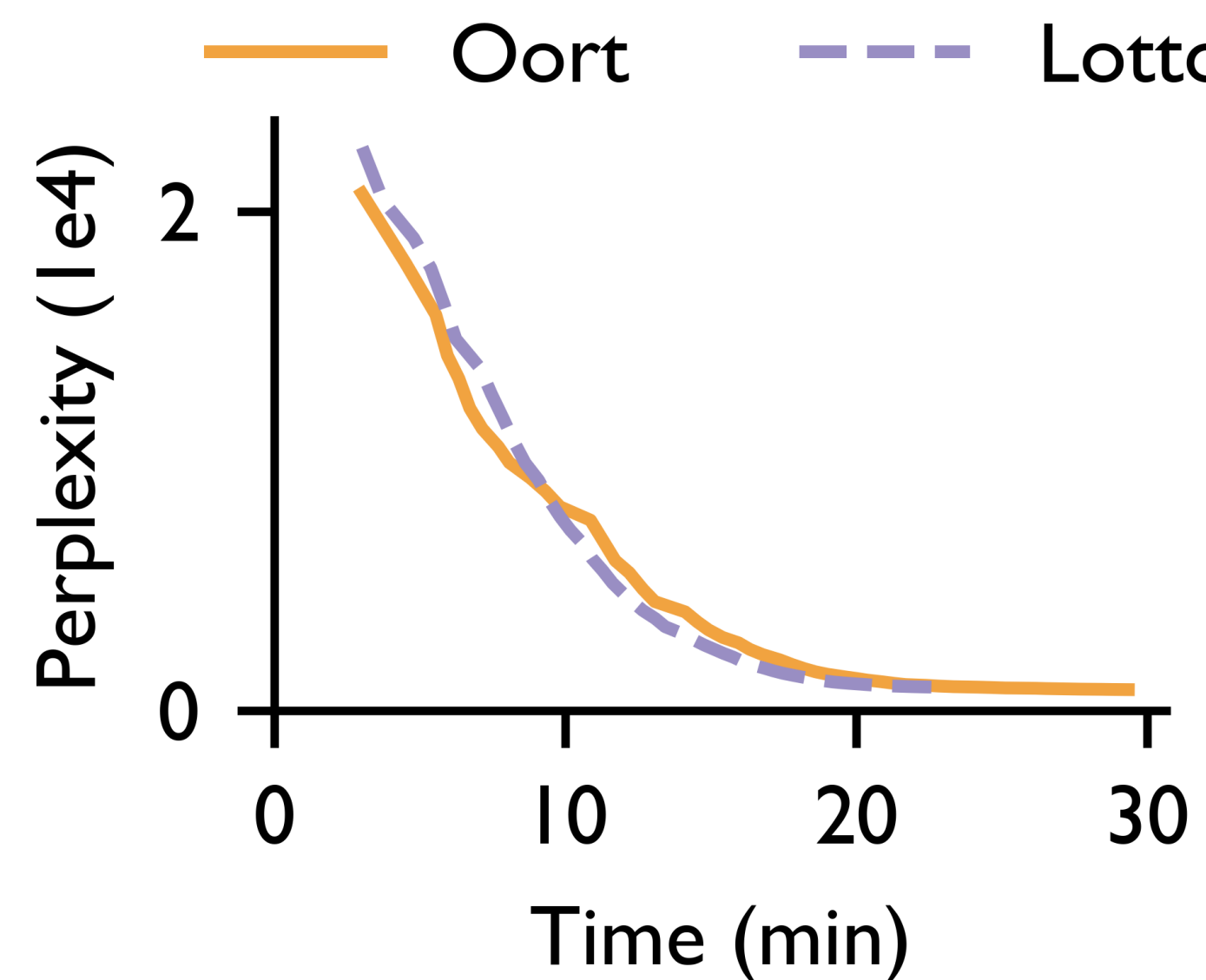
¹Lai et al. "Oort: Efficient Federated Learning via Guided Participant Selection", In OSDI '21

Lotto functions as insecure selectors

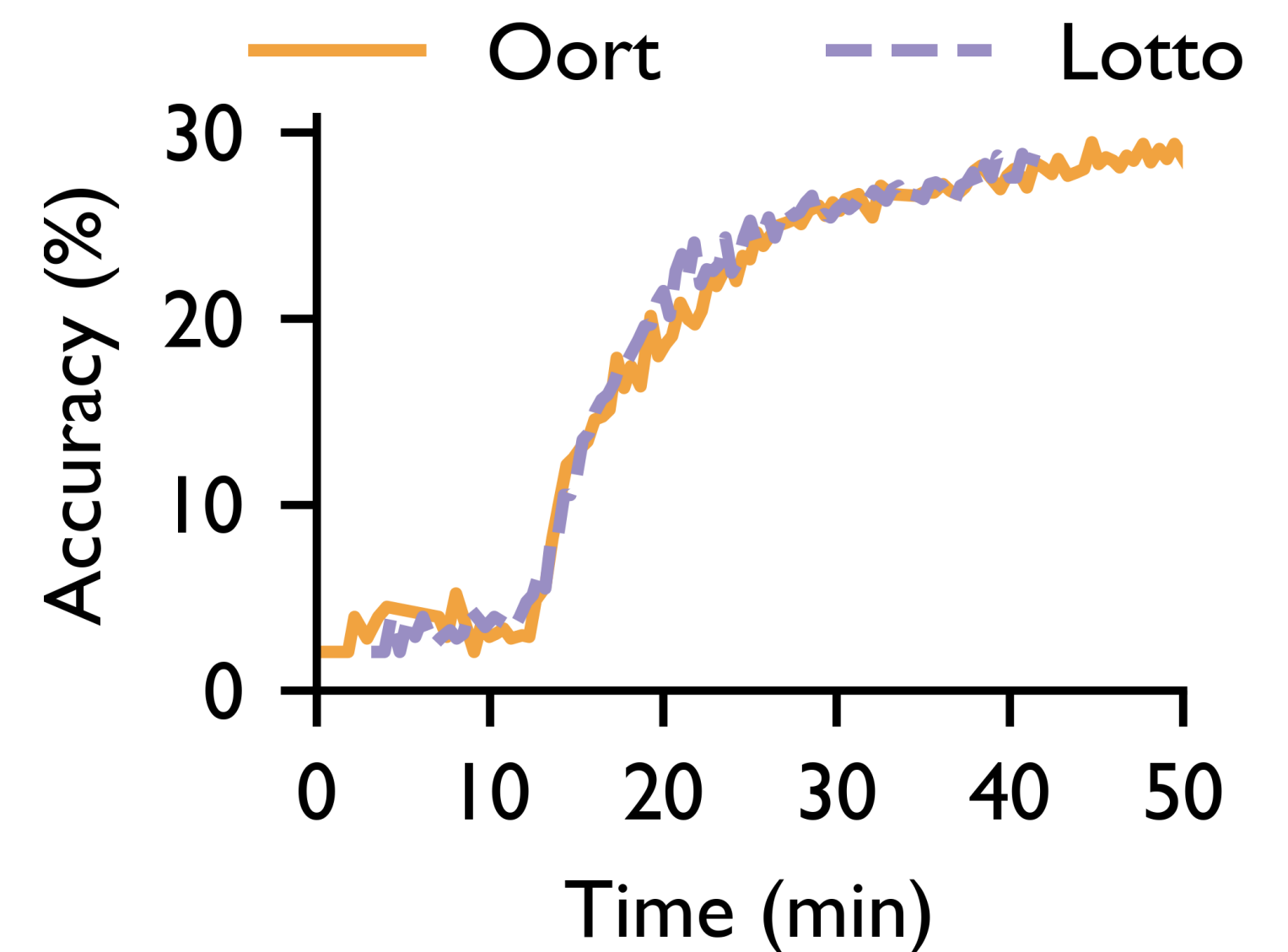
Oort¹ → State-of-the-art **informed** selector: optimized for **time-to-accuracy** of training



FEMNIST@CNN



OpenImage@MobileNet



Reddit@Albert

Lotto well approximate Oort with **no cost in time-to-accuracy** performance

¹Lai et al. "Oort: Efficient Federated Learning via Guided Participant Selection", In OSDI '21

Lotto: Results summary

Functionality

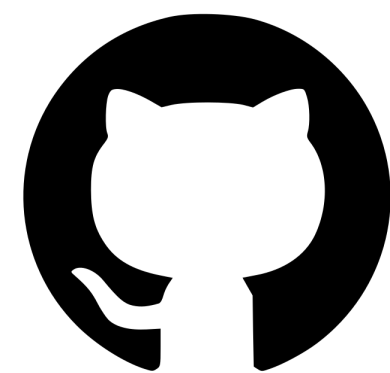
Support both **random (exact)** and **informed (well approximated)** selection

Security

Theoretical guarantee (tight probability bound) of preventing manipulation

Efficiency

Mild **runtime overhead ($\leq 10\%$)** with no **network cost ($< 1\%$)**



github.com/SamuelGong/Lotto

Thank you

zjiangaj@connect.ust.hk